

On Locally Repeated Values of Certain Arithmetic Functions, I

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Let $v(n)$ denote the number of distinct prime factors of n . We show that the equation $n + v(n) = m + v(m)$ has many solutions with $n \neq m$. We also show that if v is replaced by an arbitrary, integer-valued function f with certain properties assumed about its average order, then the equation $n + f(n) = m + f(m)$ has infinitely many solutions with $n \neq m$. © 1985 Academic Press, Inc.

1. INTRODUCTION

If $f(n)$ is an arithmetic function, one can ask for the distribution of the integers n for which $f(n) = f(n+1)$. Depending on the function f , this question is usually either trivial or intractable. However, even when the conjectured "truth" is unobtainable, partial results are sometimes possible. Also easier questions can be asked, such as: find the distribution of the n for which $|f(n) - f(n+1)|$ is small or find the distribution of the pairs n, m for which $f(n) = f(m)$ and $|n - m|$ is small.

The aim of this paper is to study the equation

$$n + v(n) = m + v(m) \quad (n \neq m) \quad (1)$$

and some related questions, where $v(n)$ is the number of distinct prime factors of n . Note that if n, m is a solution of (1), then certainly $|n - m|$ is

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small. We show that not only does (1) have many solutions, but a generalization of (1), where $\nu(n)$ is replaced by an arbitrary, integer-valued function $f(n)$ with certain properties assumed about its average order, always has infinitely many solutions.

Let $\Omega(n)$ denote the number of prime factors of n counted according to multiplicity and let $\tau(n)$ denote the number of divisors of n . Recently, D. R. Heath-Brown ("The divisor function at consecutive integers," to appear) showed that $\tau(n) = \tau(n+1)$ has infinitely many solutions. He announced that his method also gives infinitely many solutions of $\Omega(n) = \Omega(n+1)$. In a later paper in this series we shall show that the number of $n \leq x$ for which $\tau(n) = \tau(n+1)$ is $O(x/\sqrt{\log \log x})$ and the same for $\Omega(n)$ and $\nu(n)$. In another paper we shall show that $|\nu(n) - \nu(n+1)|$ is bounded on at least $cx/\sqrt{\log \log x}$ values of $n \leq x$ and the same for $\Omega(n)$. We shall also obtain an upper bound for the number of solutions of (1) and of the equation $\phi(n) = \phi(n+1)$, where ϕ is Euler's function.

Section 2 below will be devoted to the following theorem.

THEOREM 1. *Let $f(n)$ be a positive integer-valued arithmetic function for which there is a differentiable function $F(x)$ and an x_0 , such that for $x > x_0$,*

- (i)
$$\left| \sum_{n \leq x} f(n) - xF(x) \right| < \min \left\{ \frac{x}{80}, \frac{1}{240F(x)} \right\},$$
- (ii)
$$(\max_{n \leq x} f(n))^2 < \min \left\{ \frac{x}{80}, \frac{1}{240F(x)} \right\},$$
- (iii)
$$0 < F'(x) < \frac{1}{60}, \quad F'(x/2) < 3F'(x),$$
- (iv)
$$F'(x) \text{ is decreasing,}$$
- (v)
$$\lim_{x \rightarrow +\infty} F(x) = +\infty.$$

Then the equation

$$n + f(n) = m + f(m) \quad (n \neq m) \tag{2}$$

has infinitely many solutions.

COROLLARY. *Each of the equations*

$$n + \nu(n) = m + \nu(m) \quad (n \neq m)$$

$$n + \Omega(n) = m + \Omega(m) \quad (n \neq m)$$

$$n + \tau(n) = m + \tau(m) \quad (n \neq m)$$

has infinitely many solutions.

In fact, the Corollary can be derived easily from Theorem 1 by using the well-known formulas

$$\sum_{n \leq x} v(n) = x \log \log x + c_1 x + O\left(\frac{x}{\log x}\right),$$

$$\sum_{n \leq x} \Omega(n) = x \log \log x + c_2 x + O\left(\frac{x}{\log x}\right),$$

$$\sum_{n \leq x} \tau(n) = x \log x + c_3 x + O(\sqrt{x}).$$

It is not hard to construct examples to show that condition (i) on the average-order function $F(x)$ stated in Theorem 1 is nearly best possible. For example, let $f(n) = [\log \log(n+2)]$, so that it is clear that $n+f(n) = m+f(m)$ has no solutions $n \neq m$. We have

$$\sum_{n \leq x} f(n) = x \log \log x + O(x).$$

If we let $F(x) = \log \log x$, then every condition of the theorem is satisfied except that it is not true that

$$\left| \sum_{n \leq x} f(n) - x \log \log x \right| < \frac{x}{80} \quad \text{for } x > x_0.$$

Rather

$$\left| \sum_{n \leq x} f(n) - x \log \log x \right| \leq (1 + o(1)) x,$$

so that apart from the precise value of the constant our condition is best possible.

Another example is $f(n) = [n^{1/4}]$. It is again immediate that $n+f(n) = m+f(m)$ has no solutions $n \neq m$. We have

$$\sum_{n \leq x} f(n) = \frac{4}{3}x^{5/4} - \frac{1}{2}x - \left(\frac{1}{3} + 2\{x^{1/4}\}^2 - 2\{x^{1/4}\}\right)x^{3/4} + O(x^{1/2}).$$

If we let $F(x) = \frac{4}{3}x^{1/4} - \frac{1}{2} - \frac{1}{12}x^{-1/4}$, then every condition of Theorem 1 is satisfied except that

$$\left| \sum_{n \leq x} f(n) - xF(x) \right| < \frac{1}{240F'(x)} \quad \text{for } x > x_0$$

fails. Rather

$$\left| \sum_{n \leq x} f(n) - xF(x) \right| \leq \frac{1}{(20 + o(1)) F(x)},$$

so that again, apart from the precise value "1/240," the condition of the theorem is best possible.

It is interesting to note that Theorem 1 can be used to give an " Ω -theorem" for the mean value of some non-decreasing, integer-valued functions $f(n)$. Indeed, the conclusion of the theorem is not satisfied, so at least one hypothesis for the function $F(x)$ must also fail. For example, if $f(n)$ is integer-valued, non-decreasing, and $f(n) \sim \log \log n$ as $n \rightarrow \infty$, then for every value of the constant c we have

$$\limsup_{x \rightarrow \infty} \frac{1}{x} \left| \sum_{n \leq x} f(n) - x \log \log x - cx \right| \geq \frac{1}{80}.$$

In Sections 3 and 4 we shall prove (by a different method) that Eq. (1) has "many" solutions:

THEOREM 2. *There is an x_1 such that for $x > x_1$, the equation*

$$n + v(n) = m + v(m), \quad n \leq x, m \leq x, n \neq m$$

has more than $x \cdot \exp\{-4000 \log \log x \log \log \log x\}$ solutions.

2. PROOF OF THEOREM 1

Assume that the hypotheses of Theorem 1 hold. For $k = 1, 2, \dots$, let $\mathcal{S}(k)$ denote the set of integers n such that

$$0 < n \leq k \quad \text{and} \quad n + f(n) \geq k + 1.$$

Denote by $g(k)$ the number of elements of the set $\mathcal{S}(k)$.

Assume first that for some positive integer k ,

$$g(k) > g(k+1). \tag{3}$$

Clearly $k+1 \in \mathcal{S}(k+1)$ and $k+1 \notin \mathcal{S}(k)$, so that (3) implies that $\mathcal{S}(k) - \mathcal{S}(k+1)$ has at least two elements. That is, (3) implies there exist integers m, n with $m \neq n$, such that both are in $\mathcal{S}(k)$ and neither is in $\mathcal{S}(k+1)$. Then clearly, we have

$$m + f(m) = n + f(n) = k + 1,$$

giving a solution of (2). Furthermore, distinct values of k determine distinct solutions of (2). Thus if (3) holds for infinitely many k , then also (2) has infinitely many solutions.

So let us assume the opposite, that

$$g(k) \leq g(k+1) \quad \text{for } k \geq k_0. \quad (4)$$

We now show that the functions g and f have about the same average order. We have for $x > x_0$

$$\begin{aligned} \left| \sum_{k \leq x} g(k) - \sum_{n \leq x} f(n) \right| &= \left| \sum_{k \leq x} \sum_{\substack{n \leq k \\ n+f(n) \geq k+1}} 1 - \sum_{n \leq x} f(n) \right| \\ &= \left| \sum_{n \leq x} \left(-f(n) + \sum_{\substack{n \leq k < n+f(n) \\ k \leq x}} 1 \right) \right| \\ &= \left| \sum_{\substack{n \leq x \\ n+f(n) > x+1}} (-f(n) + [x] + 1 - n) \right| \\ &\leq \sum_{\substack{n \leq x \\ n > x - f(n) + 1}} f(n) < (\max_{n \leq x} f(n))^2 \\ &< \min \left\{ \frac{x}{80}, \frac{1}{240F'(x)} \right\}, \end{aligned}$$

by condition (ii) of Theorem 1. Thus by condition (i), for $x > x_0$

$$\begin{aligned} \left| \sum_{k \leq x} g(k) - xF(x) \right| &\leq \left| \sum_{k \leq x} g(k) - \sum_{n \leq x} f(n) \right| + \left| \sum_{n \leq x} f(n) - xF(x) \right| \\ &< 2 \cdot \min \left\{ \frac{x}{80}, \frac{1}{240F'(x)} \right\} = \min \left\{ \frac{x}{40}, \frac{1}{120F'(x)} \right\}. \quad (5) \end{aligned}$$

By (5) and conditions (iii) and (v) of Theorem 1 it follows that $g(k)$ is not bounded, so that for all x_1 there is an integer k_1 with

$$k_1 > x_1, g(k_1) < g(k_1 + 1). \quad (6)$$

Let x_1 be any number larger than $2 \cdot \max \{x_0, k_0, 10\}$ and let k_1 be such that (6) holds. Put

$$x = \frac{4}{3} k_1, \quad t = \left[\min \left\{ \frac{x}{4}, \frac{1}{12F'(x)} \right\} \right].$$

Then by (4) and (6) we have

$$\sum_{k=k_1+1}^{k_1+t} g(k) - \sum_{k=k_1-t+1}^{k_1} g(k) \geq t. \quad (7)$$

On the other hand, by (5), we have

$$\begin{aligned} & \sum_{k=k_1+1}^{k_1+t} g(k) - \sum_{k=k_1-t+1}^{k_1} g(k) \\ &= \sum_{k \leq k_1+t} g(k) - 2 \sum_{k \leq k_1} g(k) + \sum_{k \leq k_1-t} g(k) \\ &\leq (k_1+t)F(k_1+t) - 2k_1F(k_1) + (k_1-t)F(k_1-t) \\ &\quad + 4 \cdot \min \left\{ \frac{x}{40}, \frac{1}{120F'(x)} \right\} \\ &\leq k_1(F(k_1+t) - 2F(k_1) + F(k_1-t)) + t(F(k_1+t) - F(k_1-t)) + \frac{t}{2} \end{aligned}$$

Now condition (iv) of Theorem 1 implies F is concave downward, so that

$$\begin{aligned} & \sum_{k=k_1+1}^{k_1+t} g(k) - \sum_{k=k_1-t+1}^{k_1} g(k) < t(F(k_1+t) - F(k_1-t)) + \frac{t}{2} \\ &< 2t^2F'(k_1-t) + \frac{t}{2} \leq 2t^2F' \left(\frac{3}{4}x - \frac{1}{4}x \right) + \frac{t}{2} \\ &< 6t^2F'(x) + \frac{t}{2}, \end{aligned} \quad (8)$$

by condition (iii) of Theorem 1. Thus from (7) and (8),

$$t < 6t^2F'(x) + \frac{t}{2},$$

so that

$$t > \frac{1}{12F'(x)}. \quad (9)$$

But (9) contradicts the definition of t , which shows that there is no k_0 for which (4) holds. Thus (3) holds for infinitely many k , which, as we have seen, is sufficient for (2) to have infinitely many solutions.

3. PREPARATION FOR THEOREM 2

We shall use the following lemmas in the proof of Theorem 2.

LEMMA 1. Let \mathcal{P} be an arbitrary set of primes. Let $x \geq 1$ and let a, b be integers with $0 \leq a < b$ such that every prime factor of b is in \mathcal{P} . Let

$$g(n) = \sum_{\substack{p \leq x^{1/3}, p \notin \mathcal{P} \\ p|n}} 1, \quad E = \sum_{p \leq x^{1/3}, p \notin \mathcal{P}} \frac{1}{p}.$$

Then

$$\left| \sum_{n \leq x} (g(a + bn) - E)^2 - Ex \right| < c_4 x \tag{10}$$

where the constant c_4 is absolute (independent of each \mathcal{P}, a, b, x).

Proof. The Lemma can be proved easily by Turán's method; see [1, 2]. For the sake of completeness we give the proof.

We may assume that x is an integer. Then

$$\sum_{n \leq x} (g(a + bn) - E)^2 = \sum_{n \leq x} g(a + bn)^2 - 2E \sum_{n \leq x} g(a + bn) + xE^2. \tag{11}$$

The first sum is

$$\begin{aligned} & \sum_{n \leq x} g(a + bn)^2 \\ &= \sum_{n \leq x} \left(\sum_{\substack{p \leq x^{1/3}, p \notin \mathcal{P} \\ p|a + bn}} 1 \right)^2 \\ &= \sum_{n \leq x} \sum_{\substack{p \leq x^{1/3}, p \notin \mathcal{P} \\ p|a + bn}} 1 + \sum_{n \leq x} \sum_{\substack{p, q \leq x^{1/3}, p, q \notin \mathcal{P} \\ p \neq q, pq|a + bn}} 1 \\ &= \sum_{p \leq x^{1/3}, p \notin \mathcal{P}} \sum_{\substack{n \leq x \\ p|a + bn}} 1 + \sum_{\substack{p, q \leq x^{1/3}, p, q \notin \mathcal{P} \\ p \neq q}} \sum_{\substack{n \leq x \\ pq|a + bn}} 1 \\ &= xE + O(x^{1/3}) + x \sum_{\substack{p, q \leq x^{1/3}, p, q \notin \mathcal{P} \\ pq|a + bn}} \frac{1}{pq} + O(x^{2/3}) = xE + xE^2 + O(x), \end{aligned} \tag{12}$$

where the implied constant is absolute.

The second sum on the right of (11) is

$$-2E \sum_{n \leq x} g(a + bn) = -2E \sum_{n \leq x} \sum_{\substack{p \leq x^{1/3}, p \notin \mathcal{P} \\ p|a + bn}} 1 = -2xE^2 + O(Ex^{1/3}) \tag{13}$$

(as in the calculation in (12)) where again the implied constant is absolute.

Finally, (11), (12), and (13) yield that

$$\sum_{n \leq x} (g(a+bn) - E)^2 = xE + xE^2 - 2xE^2 + xE^2 + O(x) = xE + O(x),$$

which gives (10).

LEMMA 2. Let \mathcal{P} be a set of primes with

$$\sum_{p \in \mathcal{P}} \frac{1}{p} \leq 1. \quad (14)$$

Let a, b be integers with $0 \leq a < b$ and such that every prime factor of b is in \mathcal{P} . Let

$$f(n) = \sum_{p \notin \mathcal{P}, p|n} 1.$$

Then for

$$x > b^2, \quad (15)$$

and $x > x_2$ (where x_2 is an absolute constant, independent of \mathcal{P}, a, b) and for all $t > 0$, the number of integers n with $a + bn \leq x$ and

$$f(a+bn) < \log \log x - t\sqrt{\log \log x} \quad (16)$$

is less than $5x/t^2b$.

Proof. The lemma is clearly true if $0 < t \leq 1$, so assume now that $t > 1$. Let \mathcal{T} denote the set of integers n with $a + bn \leq x$ and satisfying (16). Define

$$g(n) = \sum_{\substack{p \leq ((x-a)/b)^{1/3}, p \notin \mathcal{P} \\ p|n}} 1,$$

so that $g(n) \leq f(n)$ for all n . Thus

$$g(a+bn) \leq f(a+bn) < \log \log x - t\sqrt{\log \log x} \quad (\text{for } n \in \mathcal{T}). \quad (17)$$

In view of (14) and (15) we have

$$E = \sum_{p \leq ((x-a)/b)^{1/3}, p \notin \mathcal{P}} \frac{1}{p} = \log \log x + O(1), \quad (18)$$

where the error term is uniformly bounded. Thus from Lemma 1 we have

$$\begin{aligned} \sum_{n \leq (x-a)/b} (g(a+bn) - E)^2 &= E \frac{x-a}{b} + O\left(\frac{x-a}{b}\right) \\ &= \frac{x}{b} (\log \log x + O(1)). \end{aligned} \quad (19)$$

On the other hand, by (17) and (18) we have for x large that

$$\begin{aligned} \sum_{n \leq (x-a)/b} (g(a+bn) - E)^2 &\geq \sum_{n \in \mathcal{F}} (g(a+bn) - E)^2 > \sum_{n \in \mathcal{F}} \left(\frac{t}{2} \sqrt{\log \log x}\right)^2 \\ &= |\mathcal{F}| \frac{t^2}{4} \log \log x. \end{aligned}$$

Combining this estimate with (19) gives the lemma.

LEMMA 3. *For $x > x_3$ and all $t > 0$, the number of integers $n \leq x$ with $v(n) \geq \log \log x + t \sqrt{\log \log x}$ is less than $5x/t^2$.*

This result is well known [1, 2] and, in fact, is a consequence of Lemma 1.

4. THE PROOF OF THEOREM 2

The idea of the proof is to show there are many disjoint intervals $[u, v] \subset [1, x]$ such that the function $n + v(n)$ maps most of $[u, v]$ into an interval $[u', v']$, where $v' - u'$ is a bit smaller than $v - u$. In fact u, v, u', v' will be found so that more integers in $[u, v]$ are mapped into $[u', v']$ than there are integers in $[u', v']$. Thus in $[u, v]$ there are at least two numbers n, m with $n + v(n) = m + v(m)$. The interval $[u, v]$ is found so that just above u , the function $v(n)$ is for most n , unusually large (so that u' can be taken large), while just below v , the function $v(n)$ behaves normally.

Let p_i denote the i th prime. Let x be a large integer, put

$$y = [42\sqrt{\log \log x}]$$

and let

$$\mathcal{P} = \{p_i : y^2 < i \leq 2y^2\},$$

$$b = \prod_{p \in \mathcal{P}} p.$$

A simple computation shows that for large enough x ,

$$\sum_{p \in \mathcal{P}} \frac{1}{p} < 1, \quad (20)$$

and

$$\begin{aligned} b &< \prod_{i \leq 2y^2} p_i = \exp((1 + o(1)) p_{2y^2}) \\ &= \exp((1 + o(1)) 2y^2 \log(2y^2)) \\ &< \exp(4000 \log \log x \times \log \log \log x) < \sqrt{x}. \end{aligned} \quad (21)$$

Let $h = h_0$ denote the least positive solution of the linear congruence system

$$h + j \equiv 0 \left(\text{mod} \prod_{i=y^2+(j-1)y+1}^{y^2+jy} p_i \right), \quad j = 1, 2, \dots, y. \quad (22)$$

Thus $0 < h_0 < b$ and h satisfies (22) if and only if

$$h \equiv h_0 \pmod{b}. \quad (23)$$

Let \mathcal{H} denote the set of integers h with $0 < h \leq x - b$ and such that (23) holds. Put

$$f(n) = \sum_{p \notin \mathcal{P}, p|n} 1.$$

For each $h \in \mathcal{H}$, let \mathcal{J}_h denote the set of positive integers j such that

$$h + j + v(h + j) < h + \log \log x + \frac{y}{2}. \quad (24)$$

(Note that if $j \in \mathcal{J}_h$, then $j < \log \log x + y/2$.) Finally, let \mathcal{H}_1 denote the set of $h \in \mathcal{H}$ with

$$|\mathcal{J}_h| < \frac{y}{16}. \quad (25)$$

We now show that at least half of the elements of \mathcal{H} are in \mathcal{H}_1 . By the construction of the sequence \mathcal{H} , for $1 \leq j \leq y$ we have

$$v(h + j) = \sum_{p|h+j} 1 = \sum_{\substack{p|h+j \\ p \in \mathcal{P}}} 1 + \sum_{\substack{p|h+j \\ p \notin \mathcal{P}}} 1 = y + f(h + j)$$

so that (24) implies

$$f(h+j) = v(h+j) - y < \log \log x - j - y/2 < \log \log x - y/2$$

(for $1 \leq j \leq y, j \in \mathcal{J}_h$). (26)

On the other hand, for $j > y$ we obtain from (24) that

$$f(h+j) \leq v(h+j) < \log \log x - j + y/2$$

(for $y < j \leq \log \log x + y/2, j \in \mathcal{J}_h$). (27)

Obviously, we have

$$\sum_{h \in \mathcal{H}} |\mathcal{J}_h| \geq \sum_{h \in \mathcal{H} - \mathcal{H}_1} \frac{y}{16} = \frac{1}{16} y |\mathcal{H} - \mathcal{H}_1|. \tag{28}$$

We now obtain an upper bound for the sum on the left of (28). We have

$$\begin{aligned} \sum_{h \in \mathcal{H}} |\mathcal{J}_h| &= \sum_{h \in \mathcal{H}} \sum_{j \in \mathcal{J}_h} 1 = \sum_{j \geq 1} \sum_{\substack{h \in \mathcal{H} \\ j \in \mathcal{J}_h}} 1 \\ &= \sum_{1 \leq j \leq y} \sum_{\substack{h \in \mathcal{H} \\ j \in \mathcal{J}_h}} 1 + \sum_{j > y} \sum_{\substack{h \in \mathcal{H} \\ j \in \mathcal{J}_h}} 1. \end{aligned}$$

By (26), the first inner sum is at most the number of terms of the arithmetic progression $h_0 + j + bn$ in $(0, x]$ with

$$f(h_0 + j + bn) < \log \log x - y/2.$$

By (27), the second inner sum is at most the number of $h_0 + j + bn$ in $(0, x]$ with

$$f(h_0 + j + bn) < \log \log x - j + y/2.$$

By (20) and (21), for large x Lemma 2 can be applied to estimate each of the inner sums, so that

$$\begin{aligned} \sum_{h \in \mathcal{H}} |\mathcal{J}_h| &\leq \sum_{1 \leq j \leq y} \frac{5x}{(y/2\sqrt{\log \log x})^2 b} + \sum_{j > y} \frac{5x}{((j-y/2)/\sqrt{\log \log x})^2 b} \\ &\leq \frac{20x \log \log x}{yb} + \frac{5x \log \log x}{b} \sum_{j > y} \frac{1}{(j-y/2)^2} \\ &< \frac{20x \log \log x}{yb} + \frac{5x \log \log x}{b} \cdot \frac{1}{y/2-1} \\ &< \frac{31x \log \log x}{yb}. \end{aligned} \tag{29}$$

Thus (28) and (29) imply that

$$|\mathcal{H} - \mathcal{H}_1| < \frac{496x \log \log x}{y^2 b} < \frac{x}{3b}.$$

With this and (21) we have for large x that

$$|\mathcal{H}_1| = |\mathcal{H}| - |\mathcal{H} - \mathcal{H}_1| > \frac{x}{b} - 2 - \frac{x}{3b} > \frac{x}{2b}, \quad (30)$$

thus showing that at least half of the elements of \mathcal{H} are in \mathcal{H}_1 . The members of \mathcal{H}_1 are our candidates for the numbers "u" described in the beginning of this section, while the numbers $h + \log \log x + y/2$ for $h \in \mathcal{H}_1$ are the candidates for the numbers "u'." However we have to do some more thinking out to allow for "v" and "v'."

Let \mathcal{H}_2 denote the set of $h \in \mathcal{H}_1$ for which there is an integer l_h with

$$h + y < l_h \leq h + b \quad (31)$$

and (letting $z = \log \log x + y/4$)

$$\sum_{\substack{h < n \leq l_h \\ n + v(n) \geq l_h + z}} 1 \leq \frac{y}{16}. \quad (32)$$

In order to give a lower bound for $|\mathcal{H}_2|$, we need an upper bound for $|\mathcal{H}_1 - \mathcal{H}_2|$, i.e., for the number of $h \in \mathcal{H}_1$ such that

$$\sum_{\substack{h < n \leq l \\ n + v(n) \geq l + z}} 1 > \frac{y}{16}$$

for all l satisfying

$$h + y < l \leq h + b.$$

We have

$$\begin{aligned} & \sum_{h \in \mathcal{H}_1 - \mathcal{H}_2} \sum_{l = h + y + 1}^{h + b} \sum_{\substack{h < n \leq l \\ n + v(n) \geq l + z}} 1 \\ & > \sum_{h \in \mathcal{H}_1 - \mathcal{H}_2} \sum_{l = h + y + 1}^{h + b} \frac{y}{16} = \sum_{h \in \mathcal{H}_1 - \mathcal{H}_2} \frac{y}{16} (b - y) \\ & > |\mathcal{H}_1 - \mathcal{H}_2| \frac{yb}{17} \end{aligned} \quad (33)$$

for large enough x .

On the other hand, by Lemma 3 we have

$$\begin{aligned}
 & \sum_{h \in \mathcal{H}_1 - \mathcal{H}_2} \sum_{l=h+y+1}^{h+b} \sum_{\substack{h < n \leq l \\ n+v(n) \geq l+z}} 1 \\
 & \leq \sum_{h \in \mathcal{H}_1 - \mathcal{H}_2} \sum_{h \leq n \leq h+b} \sum_{n \leq l \leq n+v(n)-z} 1 \\
 & \leq \sum_{n \leq x} \sum_{z \leq k \leq v(n)} 1 = \sum_{k \geq z} \sum_{\substack{n \leq x \\ v(n) \geq k}} 1 \\
 & < \sum_{k \geq z} \frac{5x}{((k - \log \log x) / \sqrt{\log \log x})^2} \\
 & = 5x \log \log x \sum_{k \geq z} \frac{1}{(k - \log \log x)^2} \\
 & < 5x \log \log x \cdot \frac{1}{y/4 - 1} < \frac{25x \log \log x}{y}
 \end{aligned} \tag{34}$$

for large x .

Thus (33) and (34) imply that

$$|\mathcal{H}_1 - \mathcal{H}_2| < \frac{17}{by} \cdot \frac{25x \log \log x}{y} = \frac{425x \log \log x}{by^2} < \frac{x}{4b},$$

so that by (30),

$$|\mathcal{H}_2| = |\mathcal{H}_1| - |\mathcal{H}_1 - \mathcal{H}_2| > \frac{x}{2b} - \frac{x}{4b} = \frac{x}{4b}. \tag{35}$$

Now for $h \in \mathcal{H}_2$ consider the interval $[h+1, l_h]$, where l_h satisfies (31) and (32). By (24) and (25), but for at most $y/16$ exceptions, every $n \in [h+1, l_h]$ has

$$n + v(n) \geq h + \log \log x + \frac{y}{2}. \tag{36}$$

By (32), but for at most $y/16$ exceptions, every $n \in [h+1, l_h]$ has

$$n + v(n) < l_h + \log \log x + \frac{y}{4}. \tag{37}$$

Thus, but for at most $y/8$ exceptions, every $n \in [h+1, l_h]$ has both (36) and (37) holding. So we have at least $l_h - h - y/8$ numbers n mapped by $n + v(n)$ to the interval

$$\left[h + \log \log x + \frac{y}{2}, l_h + \log \log x + \frac{y}{4} \right),$$

which has at most $l_h - h - \lfloor y/4 \rfloor$ integers. There are therefore at least $\lfloor y/8 \rfloor$ pairs $n, m \in [h+1, l_h]$ with $n \neq m$ and $n + v(n) = m + v(m)$.

As h runs over \mathcal{H}_2 , the intervals $[h+1, l_h]$ are disjoint and contained in $[1, x]$. Thus, below x , there are at least (using (21), (35) and assuming x is large)

$$\lfloor y/8 \rfloor |\mathcal{H}_2| > 4 |\mathcal{H}_2| > \frac{x}{b} > x \cdot \exp(-4000 \log \log x \log \log \log x)$$

pairs $n, m \leq x$ with $n \neq m$ and $n + v(n) = m + v(m)$.

Remarks. Probably Theorem 2 is far from the truth. We conjecture that there are positive constants c_5, c_6 such that

$$|\{n \leq x : \exists m \neq n, n + v(n) = m + v(m)\}| \sim c_5 x,$$

$$|\{(n, m) : n \leq x, m \leq x, n + v(n) = m + v(m)\}| \sim c_6 x.$$

Almost the same proof as for Theorem 2 can show the analogous result with $\Omega(n)$ replacing $v(n)$. With a little more difficulty, the same can be proved with $\tau(n)$ replacing $v(n)$.

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