

DENOMINATORS OF EGYPTIAN FRACTIONS II

BY

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I. Introduction

A positive fraction a/N is said to be written in Egyptian form if we write

$$a/N = 1/n_1 + 1/n_2 + \cdots + 1/n_k, \quad 0 < n_1 < n_2 < \cdots < n_k,$$

where the n_i are integers. Among the many expansions for each fraction a/N there is some expansion for which n_k is minimal. Let $D(a, N)$ denote the minimal value of n_k .

Define $D(N)$ by $D(N) = \max \{D(a, N) : 0 < a < N\}$. We are interested in the behavior of $D(N)$. In our paper [1] we showed that for $N = P$, a prime, $D(P) \geq P \log P$ and that for some constant K and any $N > 1$, $D(N) \leq KN (\log N)^4$. It was surprising that such close upper and lower bounds could be achieved by the simple techniques of [1]. In this paper we refine the techniques of [1] and show that on the one hand for P large enough that $\log_{2r} P \geq 1$,

$$D(P) \geq \frac{P \log P \log_2 P}{\log_{r+1} P \prod_{j=4}^{r+1} \log_j P}$$

and on the other hand that for $\varepsilon > 0$ and N sufficiently large (Theorem 1 and its corollary yield more precise statements), $D(N) \leq (1 + \varepsilon)N (\log N)^2$. We conjecture that the exponent 2 can be replaced by $(1 + \delta)$ for $\delta > 0$.

As part of the proof of the above results we need to analyze the number of distinct subsums of the series $\sum_{i=1}^N 1/i$, say $S(N)$. We show that whenever $\log_{2r} N \geq 1$,

$$\frac{\alpha N}{\log N} \prod_{j=3}^r \log_j N \leq \log S(N) \leq \frac{N \log_r N}{\log N} \prod_{j=3}^r \log_j N$$

for some $\alpha \geq 1/e$.

II. The upper bound for $D(N)$

Let p_k denote the k th prime, and let $\Pi_k = \prod_{i=1}^k p_i$. We recall from [1]:

LEMMA 1. *If $0 < r < \sigma(\Pi_k)$ then there are divisors d_i of Π_k such that $r = \sum d_i$.*

Received July 5, 1974; received in revised form January 27, 1976.

LEMMA 2. For N sufficiently large, if k is chosen so that $\Pi_{k-1} \leq N \leq \Pi_k$, then

$$p_k \leq \log N \left(1 + \frac{2}{\log \log N} \right).$$

Proof. If $\vartheta(x) = \sum_{p \leq x} \log p$ then $\log \Pi_k = \vartheta(p_k)$. We note that p_k is the least prime such that $\vartheta(p_k) \geq \log N$. By [4, Theorem 4], $\vartheta(x) \geq x(1 - (1/2 \log x))$ for large enough x . Thus if

$$x_0 = \log N \left(1 + \frac{1}{\log \log N} \right)$$

then $\vartheta(x_0) \geq \log N$. Let p_0 be the least prime greater than x_0 . For x_0 sufficiently large we have [3, p. 323] $p_0 \leq x_0 + x_0^{2/3}$. Since $p_k \leq p_0$,

$$p_k \leq \log N \left(1 + \frac{2}{\log \log N} \right)$$

for N sufficiently large.

LEMMA 2*. If $N \geq 2$ and $\Pi_{k-1} < N \leq \Pi_k$ then $p_k \leq 2 \log N / \log 2$.

Proof. For $N = 2$, $p_k = 2$ and the lemma holds. For $3 \leq N \leq 6$, $p_k = 3$ and the lemma holds. For $\Pi_2 < N \leq \Pi_{16}$ the theorem follows since for $k \leq 16$, computation shows that $p_k \leq 2 \log \Pi_{k-1} / \log 2$. For $N \geq \Pi_{16}$ we have $\log N \geq 41$. By definition of $\vartheta(x)$, $\log \Pi_k = \vartheta(p_k)$ where p_k is the least prime such that $\vartheta(p_k) \geq \log N$. Since for $x \geq 41$ we have [4, Theorem 4, Corollary] $\vartheta(x) \geq x(1 - (1/\log x))$, we see that

$$\vartheta(x_0) \geq \log N \quad \text{for } x_0 = \log N \left(1 + \frac{3}{2 \log \log N} \right) \geq 41.$$

By Bertrand's postulate we see that $p_k \leq 2x_0$. Since

$$2 \left(1 + \frac{3}{2 \log \log N} \right) \leq 2 / \log 2 \quad \text{when } \log N \geq 41,$$

the lemma follows.

LEMMA 3. If $N \geq 12$, then in the closed interval $[\sqrt{N}, N + \sqrt{N}]$ there are at least $[N/2] + 1$ square-free integers with all prime factors less than N .

Proof. Let $\Pi^* = \prod_{p < N} p$. Let $D = \{m: \sqrt{N} \leq m \leq N + \sqrt{N}, m \mid \Pi^*\}$. Let $Q(x)$ be the number of square free integers not exceeding x . Thus

$$|D| \geq Q(N + \sqrt{N}) - Q(\sqrt{N}) - L$$

where L is the number of primes between N and $N + \sqrt{N}$ inclusive. Suppose $N \geq 24^2$, so that $\sqrt{N} \geq 24$. In the interval $[N, N + \sqrt{N}]$ only odd numbers can be prime; there are at most $1 + \frac{1}{2} \sqrt{N}$ odd numbers, and at least four of

them are divisible by 3. We deduce that $L \leq (\frac{1}{2}\sqrt{N}) - 3$. From the proof of Theorem 333 in [2] we see that

$$Q(x) = \sum_{d^2 \leq x} \mu(d) \left[\frac{x}{d^2} \right].$$

Thus

$$\begin{aligned} Q(N + \sqrt{N}) - Q(\sqrt{N}) &= \sum_{d \leq \sqrt{N + \sqrt{N}}} \mu(d) \left[\frac{N + \sqrt{N}}{d^2} \right] \\ &\quad - \sum_{d \leq N^{1/4}} \mu(d) \left[\frac{\sqrt{N}}{d^2} \right] \\ &\geq (N + \sqrt{N}) \sum_{d \leq \sqrt{N + \sqrt{N}}} \frac{\mu(d)}{d^2} \\ &\quad - \sqrt{N} \sum_{\substack{d \leq N^{1/4} \\ d \leq N^{1/4}}} \frac{\mu(d)}{d^2} - [\sqrt{N + \sqrt{N}}]. \end{aligned}$$

Since $\sum_{d=1}^{\infty} \mu(d)/d^2 = 1/\zeta(2) = 6/\pi^2$ and $|\mu(d)| \leq 1$ we get

$$\begin{aligned} Q(N + \sqrt{N}) - Q(\sqrt{N}) &\geq \frac{6N}{\pi^2} - [\sqrt{N + \sqrt{N}}] - N \sum_{d > \sqrt{N + \sqrt{N}}} \frac{1}{d^2} \\ &\quad - \sqrt{N} \sum_{N^{1/4} < d \leq \sqrt{N + \sqrt{N}}} \frac{1}{d^2} \\ &> \frac{6N}{\pi^2} - M - \frac{N}{M} - \sqrt{N} \left(\frac{1}{[N^{1/4}]} - \frac{1}{M} \right) \end{aligned}$$

where $M = [\sqrt{N + \sqrt{N}}]$. Since $\sqrt{N + \sqrt{N}} - \sqrt{N} - \sqrt{N} \geq 1$, we see that $M \geq \sqrt{N} - \sqrt{N}$ and hence that the above expression is decreasing in M . Thus we obtain

$$\begin{aligned} Q(N + \sqrt{N}) - Q(\sqrt{N}) &\geq \frac{6N}{\pi^2} - \sqrt{N + \sqrt{N}} - \frac{N}{\sqrt{N + \sqrt{N}}} \\ &\quad - \sqrt{N} \left(\frac{1}{[N^{1/4}]} - \frac{1}{\sqrt{N + \sqrt{N}}} \right) \\ &= \frac{6N}{\pi^2} - \frac{2N}{\sqrt{N + \sqrt{N}}} - \frac{\sqrt{N}}{[N^{1/4}]}. \end{aligned}$$

Thus

$$|D| \geq \frac{6N}{\pi^2} - \frac{2N}{\sqrt{N + \sqrt{N}}} - \frac{\sqrt{N}}{[N^{1/4}]} - \frac{\sqrt{N}}{2} + 3.$$

To show that $|D| \geq N/2$ it suffices to show that

$$0.1079 \cdots = \frac{6}{\pi^2} - \frac{1}{2} \geq \frac{2}{\sqrt{N} + \sqrt{N}} + \frac{1}{2\sqrt{N}} + \frac{1}{\sqrt{N[N^{1/4}]}} - \frac{3}{N}$$

which is true for $N = 24^2$, whence for $N \geq 24^2$. On the other hand one can verify directly and/or by special arguments that the lemma is true for $576 \geq N \geq 12$.

LEMMA 4. If $\Pi_k(1 - (2/\sqrt{p_k})) \leq r < 2\Pi_k$ then there are distinct d_i such that $d_i | \Pi_k$, $d_i > \Pi_{k-1}(p_k + \sqrt{p_k})^{-1}$ and $r = \sum d_i$.

Proof. We note, in order to begin a proof by induction, that the lemma is true for $k = 1, 2, 3$, since for these cases $\Pi_{k-1}(p_k + \sqrt{p_k})^{-1} < 1$. We suppose $k \geq 4$ and that the lemma is true for all $k' < k$. Consider the set

$$D = \{d: \sqrt{p_k} \leq d < p_k + \sqrt{p_k}, d | \Pi_{k-1}\}.$$

Case 1. $k \geq 6$, i.e., $p_k \geq 13$. Let r be given in the desired range. According to Lemma 3, $|D| \geq (p_k + 1)/2$. Also note that no two elements of D are congruent mod p_k and that none is congruent to zero mod p_k . Let

$$D^* = \{0\} \cup \{\Pi_{k-1}/d; d \in D\}.$$

If $d \in D^*$, $d \neq 0$ then $\Pi_{k-1}(\sqrt{p_k})^{-1} \geq d \geq \Pi_{k-1}(p_k + \sqrt{p_k})^{-1}$. We note that $|D^*| \geq (p_k + 3)/2$ and no two elements of D^* are congruent mod p_k . If $r \equiv 2d \pmod{p_k}$ for some $d \in D^*$, let $D^{**} = D^* \setminus \{d\}$, otherwise let $D^{**} = D^*$. Hence $|D^{**}| \geq (p_k + 1)/2$ and we may apply the Cauchy-Davenport Theorem to find d' and d'' , distinct elements of D^{**} such that $r - d' - d'' \equiv 0 \pmod{p_k}$. Let $r^* = r - d' - d''$. Then

$$r^* \geq r - \frac{2\Pi_{k-1}}{\sqrt{p_k}} \geq \Pi_k \left(1 - \frac{2}{\sqrt{p_k}} - \frac{2}{p_k\sqrt{p_k}}\right).$$

Since $1/\sqrt{p_{k-1}} - 1/\sqrt{p_k} \geq 1/p_k\sqrt{p_k}$, as is seen by using the mean value theorem on $1/\sqrt{x}$, we deduce that $r^* \geq \Pi_k(1 - (2/\sqrt{p_{k-1}}))$. Let $r' = r^*/p_k$, an integer. Then

$$\Pi_{k-1} \left(1 - \frac{2}{\sqrt{p_{k-1}}}\right) \leq r' < 2\Pi_{k-1},$$

so by induction $r' = \sum d_i$ where $d_i | \Pi_{k-1}$, $d_i \geq (p_{k-1} + \sqrt{p_{k-1}})^{-1} \Pi_{k-2}$. It follows that $r = \sum p_k d_i + d' + d''$, and since the d_i were distinct by induction, so are the $p_k d_i$; also, unless either d' or d'' is zero, in which case we discard it from the sum, $d', d'' \not\equiv 0 \pmod{p_k}$ so that all the terms in the sum are distinct. Clearly

$$d', d'' \geq \frac{\Pi_{k-1}}{p_k + \sqrt{p_k}}.$$

On the other hand, by induction

$$d_i \geq \frac{\Pi_{k-2}}{p_{k-1} + \sqrt{p_{k-1}}},$$

thus

$$d_i p_k \geq \frac{\Pi_{k-2} p_k}{p_{k-1} + \sqrt{p_{k-1}}} \geq \frac{\Pi_{k-1}}{p_k + \sqrt{p_k}}.$$

Case 2. $k = 4, 5, p_k = 7, 11$. An easy computation shows that for $p_k = 7$, $D^* = \{0, 5, 6, 10\}$. Every nonzero congruence class mod 7 can be obtained as a sum of two or fewer elements of D^* as follows: $1 \equiv 5 + 10$, $2 \equiv 6 + 10$, $3 \equiv 10 + 0$, $4 \equiv 5 + 6$, $5 \equiv 5 + 0$, and $6 \equiv 6 \pmod{7}$. Thus for $r \not\equiv 0 \pmod{7}$ we may proceed to define r' as in Case 1. If $r \equiv 0 \pmod{7}$, let $r^* = r$ and proceed as in Case 1.

For $p_k = 11$, $D^* = \{0, 2 \cdot 3 \cdot 7, 5 \cdot 7, 2 \cdot 3 \cdot 5, 3 \cdot 7, 3 \cdot 5\} \equiv \{0, 9, 2, 8, 10, 4\} \pmod{11}$. Every congruence class mod 11 can be obtained as a sum of at most three distinct elements of D^* as follows: $0 \equiv 0$, $1 \equiv 10 + 2$, $2 \equiv 2$, $3 \equiv 10 + 4$, $4 \equiv 4$, $5 \equiv 10 + 4 + 2$, $6 \equiv 4 + 2$, $7 \equiv 10 + 8$, $8 \equiv 10 + 9$, $9 \equiv 9$, $10 \equiv 10$. Thus we may define r' and proceed as in Case 1. The proof is completed.

We are now ready to prove:

THEOREM 1. For every N , $D(N) \leq \lambda^3(N)N(\ln N)^2$ where $2/\log 2 \geq \lambda(N) \geq 1$ and $\lim_{N \rightarrow \infty} \lambda(N) = 1$.

Proof. Given a/N choose Π_k such that $\Pi_{k-1} < N \leq \Pi_k$. If $N \mid \Pi_k$, then $a/N = b/\Pi_k$. By Lemma 1, $b = \sum d_i$, $d_i \mid \Pi_k$. By reducing the fractions in $\sum d_i/\Pi_k$ we obtain a representation of a/N in which no denominator exceeds $\Pi_k < 2N \log N/\log 2$.

If $N \nmid \Pi_k$ write $a/N = (qN + r)/N\Pi_k$ where r is chosen so that

$$\Pi_k \left(1 - \frac{2}{\sqrt{p_k}}\right) \leq r \leq 2\Pi_k.$$

This can be done since we may assume $a \geq 2$ and since $N \leq \Pi_k$. The fraction q/Π_k can be handled by Lemma 1, as in the paragraph above. We now use Lemma 4 to write r/Π_k in Egyptian form using very small denominators. By Lemma 4, $r = \sum d_i$ where $d_i \mid \Pi_k$, the d_i are distinct and $d_i \geq \Pi_{k-1}(p_k + \sqrt{p_k})^{-1}$. Thus $r/\Pi_k = (\sum d_i)/\Pi_k = \sum 1/n'_i$ where $n'_i = \Pi_k/d_i$. Thus the n'_i are distinct and $n'_i \leq p_k(p_k + \sqrt{p_k})$. It follows that $r/N\Pi_k = \sum 1/n_i$ where $n_i = n'_i N$ and the n_i are distinct from each other as well as from the denominators in the expansion of q/Π_k since these denominators all divide Π_k while $N \mid n_i$ and $N \nmid \Pi_k$. Furthermore

$$n_i \leq N p_k (p_k + \sqrt{p_k}) \leq \lambda^3(N)N(\ln N)^2$$

where $\lambda(N)$ can be chosen to satisfy $2/\log 2 \geq \lambda(N)$ by Lemma 2*, $\lim_{N \rightarrow \infty} \lambda(N) = 1$ by Lemma 2, and $\lambda(N) \geq (1 + (1/\sqrt{\log N}))$.

III. The number of distinct subsums of $\sum_{i=1}^N 1/i$.

DEFINITION. Let $S(N)$ denote the number of distinct values of $\sum_{k=1}^N \varepsilon_k/k$ where the ε_k 's take on all possible combinations of values with $\varepsilon_k = 0$ or 1.

To obtain a lower bound for $S(N)$ we begin with the following lemma.

LEMMA 5. For all $N \geq 3$, $S(N) \geq 2^{N/\log N}$.

Proof. It is clear that each distinct choice of the ε_p 's for p prime yields a different value of $\sum_{p \leq N} \varepsilon_p/p$. Thus $S(N) \geq 2^{\pi(N)}$. Since for $N \geq 17$, $\pi(N) \geq N/\log N$ by Corollary 1 of Theorem 2 of [4], the lemma is true for $N \geq 17$. To verify that the result holds for $3 \leq N \leq 16$, note that both $S(N)$ and $2^{N/\log N}$ are monotone and $2^{4/\log 4} \leq 8 \leq S(3)$, $2^{12/\log 12} < 2^5 \leq S(5)$ and $2^{16/\log 16} < 2^6 = 2^{\pi(13)} \leq S(13)$, where $S(3) = 8$ and $S(5) = 2^5$ are a result of direct verification. Thus the lemma is proved.

THEOREM 2. If $r \geq 1$ and N is large enough that $\log_{2^r} N \geq 1$, then

$$S(N) \geq \exp\left(\alpha \cdot \frac{N}{\log N} \cdot \prod_{j=3}^r \log_j N\right)$$

where $\alpha = 1/e$ is a permissible value for α and $\log_1 x = \log x$, $\log_j x = \log(\log_{j-1} x)$.

Proof. The proof is by induction on r .

In order to prove the theorem with the proper constant we make the slightly stronger (as will be shown at the end of the proof) inductive hypothesis

$$(*) \quad S(N) \geq \exp\left(\prod_{j=3}^k \left(1 - \frac{3}{\log_{2^{j-2}} N}\right) \cdot \frac{N}{\log N} \prod_{j=3}^k \log_j N\right)$$

for $\log_{2^k} N \geq 1$. The hypothesis (*) is clearly true for $k = 1, 2$ by Lemma 5. We assume the induction hypothesis holds for $k = 1, 2, \dots, r-1$ and show that it also holds for $k = r \geq 3$.

Let $Q = 2N/\log N$ and $Q' = N/\log_2 N$. Note that $Q' > Q$. We define \mathscr{P} by

$$\mathscr{P} = \{N \geq p \geq Q: p \text{ a prime}\}.$$

Let $T = \{k \leq N: \text{there exists } p \in \mathscr{P}, p \mid k\}$.

$S(N)$ is greater than the number of distinct values of the sum $\sum_{k \in T} \varepsilon_k/k$, which we denote by $T(N)$. We rewrite the sum as

$$\sum_{k \in T} \frac{\varepsilon_k}{k} = \sum_{p \in \mathscr{P}} \frac{1}{p} \left(\sum_{k=1}^{N/p} \frac{\varepsilon_k}{k} \right).$$

Set $\sum_{k=1}^{N/p} \varepsilon_k/k = a_p/b_p$ where $\log b_p = \psi(N/p)$, $\psi(x) = \sum_{p \leq x} \log p$. Also

$$a_p \leq 2b_p \log N/p \quad \text{for } p \leq N/3.$$

Thus, if

$$\frac{1}{p} \left(\frac{a_p}{b_p} - \frac{a'_p}{b'_p} \right) = \frac{c}{d}, \quad (c, d) = 1,$$

then $p \mid d$ if $p \nmid (a_p - a'_p)$. But for $p \leq N/3$,

$$a_p - a'_p \leq 2b_p \log N/p \leq 2 \log(N/p) e^{\psi(N/p)}.$$

Since $\psi(x) < (1.04)x$ [4, Theorem 12] we see that

$$a_p - a'_p \leq 2 \log(N/p) e^{(1.04)N/p} < Q \leq p,$$

since $N \geq e^e$. For $p > N/3$ it is clear that $p \nmid (a_p - a'_p)$.

Thus $p \nmid (a_p - a'_p)$ and $p \mid d$. It follows that distinct choices of a_p/b_p yield distinct sums. Thus $T(N) \geq \prod_{p \in \mathcal{P}} S(N/p)$, so that $S(N) \geq \prod_{p \in \mathcal{P}} S(N/p)$.

We will now evaluate the above product using our inductive hypothesis. First note that

$$\log S(N) \geq \sum_{p \in \mathcal{P}} \log S\left(\frac{N}{p}\right).$$

For simplicity let $S^*(x) = \log S(x)$.

We recall the well-known method using Stieltjes integration with respect to $\vartheta(x)$ and integration by parts by which one evaluates sums where the variable runs over primes [4, p. 74].

LEMMA 6. *If $f'(p)$ exists and is continuous then*

$$\begin{aligned} \sum_{Q < p \leq Q'} f(p) &= \int_Q^{Q'} \frac{f(x)}{\log x} dx + \left(\frac{\vartheta(x) - x}{\log x} f(x) \right) \Big|_Q^{Q'} \\ &\quad - \int_Q^{Q'} (\vartheta(x) - x) \frac{d}{dx} \left(\frac{f(x)}{\log x} \right) dx. \end{aligned}$$

Let $L^*(x) = x/\log x \prod_{j=3}^{r-1} \log_j x$, and note that for $Q < p \leq Q'$, $N/p \geq \log_2 N$; hence $\log_{2(r-1)} N/p \geq \log_{2r} N \geq 1$, and the induction assumption tells us that

$$S^*(N/p) \geq \prod_4^r \left(1 - \frac{3}{\log_{2j-2} N} \right) L^*(N/p).$$

We thus obtain

$$\begin{aligned} \left(\prod_{j=4}^r \left(1 - \frac{3}{\log_{2j-2} N} \right) \right)^{-1} S^*(N) &\geq \sum_{Q < p \leq Q'} L^*(N/p) \\ &= \int_Q^{Q'} \frac{L^*(N/x)}{\log x} dx + \frac{\vartheta(x) - x}{\log x} L^*(N/x) \Big|_Q^{Q'} \\ &\quad - \int_Q^{Q'} (\vartheta(x) - x) \frac{d}{dx} \left(\frac{L^*(N/x)}{\log x} \right) dx \\ &= S_1 + S_2 + S_3, \quad \text{say.} \end{aligned}$$

We shall estimate the absolute values of S_2 and S_3 and then the value of S_1 , the main term. We use the estimate [4, p. 70] $|\theta(x) - x| < x/(2 \log x)$ to obtain

$$|S_2| \leq \frac{N}{\log^2 N} \prod_4^r \log_j N$$

as follows:

$$\begin{aligned} |S_2| &\leq \frac{Q'}{2 \log^2 Q'} L^*(N/Q') + \frac{Q}{2 \log^2 Q} L^*(N/Q) \\ &= \frac{NL^*(\log_2 N)}{2 \log_2 N (\log N - \log_3 N)^2} + \frac{NL^*\left(\frac{\log N}{2}\right)}{\log N (\log N + \log 2 - \log_2 N)^2} \\ &\leq \frac{N}{2 \log^2 N \cdot \log_2 N \left(1 - \frac{\log_3 N}{\log N}\right)^2} \cdot \frac{\log_2 N}{\log_3 N} \prod_5^{r+1} \log_j N \\ &\quad + \frac{N}{\log^3 N \left(1 + \frac{\log 2 - \log_2 N}{\log N}\right)^2} \cdot \frac{\log N}{2 (\log_2 N - \log 2)} \prod_4^r \log_j N \\ &\leq \frac{N}{2 \log^2 N} \prod_4^r \log_j N \\ &\quad \cdot \left(\frac{\log_{r+1} N}{\log_3 N \log_4 N \left(1 - \frac{\log_3 N}{\log N}\right)^2} + \frac{1}{\left(1 - \frac{\log_2 N}{\log N}\right)^2 (\log_2 N - \log 2)} \right) \\ &\leq \frac{N}{\log^2 N} \prod_4^r \log_j N. \end{aligned}$$

A straightforward calculation yields

$$\left| \frac{d}{dx} \frac{L^*(N/x)}{\log x} \right| \leq \frac{N}{x^2 \log x \log N/x} \prod_3^{r-1} \log_j N/x$$

for x in the prescribed range. Thus

$$|S_3| \leq \int_Q^{Q'} \frac{N}{2x \log^2 x \log N/x} \prod_3^{r-1} \log_j N/x \, dx.$$

Using the facts that $N/x \leq \log N$ and $2 \log^2 x \geq (3/2) \log^2 N$ for all x in the range of integration, we see that

$$\begin{aligned} |S_3| &\leq \frac{2N \prod_4^r \log_j N}{3 \log^2 N} \int_Q^{Q'} \frac{dx}{x \log N/x} \\ &= \frac{2N \prod_4^r \log_j N}{3 \log^2 N} (-\log_2 N/x|_Q^{Q'}) \\ &\leq \frac{2N \prod_4^r \log_j N}{3 \log^2 N} (-\log_2 N/x|_{N/\log N}^{Q'}) \\ &\leq \frac{N \prod_3^r \log_j N}{\log^2 N}. \end{aligned}$$

We next obtain a lower bound for S_1 :

$$\begin{aligned} S_1 &= \int_Q^{Q'} \frac{N}{x \log x \log N/x} \prod_3^{r-1} \log_j N/x \, dx \\ &\geq \frac{N}{\log N} \int_Q^{Q'} \frac{\prod_3^{r-1} \log_j N/x}{x \log N/x} \, dx. \end{aligned}$$

With $u = \prod_3^{r-1} \log_j N/x$ and $v = -\log_2 N/x$ we integrate by parts to obtain

$$\begin{aligned} &\int_Q^{Q'} \frac{1}{x \log N/x} \cdot \prod_3^{r-1} \log_j N/x \, dx \\ &= - \prod_2^{r-1} \log_j N/x|_Q^{Q'} - \int_Q^{Q'} \frac{1}{x \log N/x} \left(\sum_{i=3}^{r-1} \prod_{j=i+1}^{r-1} \log_j N/x \right) dx \\ &\geq \prod_2^{r-1} \log_j N/Q - \prod_4^{r+1} \log_j N - 2 \prod_5^r \log_j N/Q \int_Q^{Q'} \frac{dx}{x \log N/x} \\ &\geq \prod_3^r \log_j N \left(1 - \frac{5}{2 \log_4 N} \right), \end{aligned}$$

where we have used that

$$\prod_2^{r-1} \log_j \frac{x}{2} \geq \left(1 - \frac{2}{\log x} \right) \prod_2^{r-1} \log_j x \quad \text{for } \log_2 N \leq x \leq \log N.$$

Substituting this in the lower bound for S_1 we obtain

$$S_1 \geq \frac{N}{\log N} \cdot \prod_3^r \log_j N \left(1 - \frac{5}{2 \log_4 N} \right).$$

Combining the estimates for $S_1, |S_2|$ and $|S_3|$ we obtain

$$\begin{aligned} & \left(\prod_{j=4}^r \left(1 - \frac{3}{\log_{2^{j-2}} N} \right) \right)^{-1} S^*(N) \\ & \geq \frac{N}{\log N} \prod_3^r \log_j N \left\{ 1 - \frac{5}{2 \log_4 N} - \frac{2}{\log N \log_3 N} - \frac{1}{\log N} \right\} \\ & \geq \left(1 - \frac{3}{\log_4 N} \right) \frac{N}{\log N} \prod_3^r \log_j N \end{aligned}$$

which satisfies (*). Thus (*) holds for all $r \geq 1$.

Since we know $\log_{2^r} N \geq 1$ we deduce that $\log_{2^{j-2}} N \geq e^{2r-2j+2}$. Thus

$$\begin{aligned} \prod_{j=3}^r \left(1 - \frac{3}{\log_{2^{j-2}} N} \right) & \geq \prod_{j=3}^r \left(1 - \frac{3}{e^{2r-2j+2}} \right) \\ & = \prod_{j=1}^{r-2} \left(1 - \frac{3}{e^{2j}} \right) \\ & \geq \prod_1^\infty \left(1 - \frac{3}{e^{2j}} \right) \\ & \geq 1/e, \end{aligned}$$

where the last inequality follows from the facts that for $0 \leq x \leq 3/e^2 = 0.406 \dots$, $\log(1-x) \geq -3x/2$ and $-(3/2) \sum_{j=1}^\infty 3/e^{2j} = -0.526 \dots > -1$.

The theorem is proved.

LEMMA 7. For $N \geq 1, S(N) \leq 2^N$.

Proof. The result follows immediately since there are 2^N distinct choices for $e_i, 1 \leq i \leq N, e_i = 0$ or 1 .

LEMMA 8. For $\log_2 N \geq 1, S(N) \leq \exp(N/\log_2 N)$.

For $\log_4 N \geq 1, S(N) \leq \exp(N \log_2 N / \log N)$.

Proof of Lemma 8. Let $Q = N/\log N$. Let

$$\mathcal{P} = \{p: Q < p \leq N\},$$

$$Z_1 = \{k \leq N: \text{there exists } p \in \mathcal{P}, p \mid k\}$$

and

$$Z_2 = \{k \leq N: k \notin Z_1\}.$$

Thus we may write

$$\sum_{k=1}^N \frac{e_k}{k} = \sum_{k \in Z_1} \frac{e_k}{k} + \sum_{k \in Z_2} \frac{e_k}{k}.$$

Let $S_i(N)$ denote the number of distinct values of the sum with $k \in Z_i$ as the e_k 's take on all possible values with $e_k = 0$ or 1 . As before $S^*(N) = \log S(N)$ and $S_i^*(N) = \log S_i(N), i = 1, 2$.

The case $\log_2 N \geq 1$.

Subcase A. $N \geq 10^8$. We estimate $S_1^*(N)$ first. From the definition of Z_1 we see that

$$|Z_1| = \sum_{p \in \mathcal{P}} \left[\frac{N}{p} \right] \leq N \sum_{p \in \mathcal{P}} \frac{1}{p}.$$

Using the estimates of [4, Theorem 5 and corollary], we obtain

$$|Z_1| \leq N \left(\log_2 N - \log_2 Q + \frac{1}{\log^2 N} + \frac{1}{2 \log^2 Q} \right).$$

Since $S_1(N) \leq 2^{|Z_1|}$, it follows that

$$(1) \quad S_1^*(N) \leq N (\log 2) \left(\log_2 N - \log_2 Q + \frac{1}{\log^2 N} + \frac{1}{2 \log^2 Q} \right).$$

We now estimate $S_2(N)$. Suppose $\sum_{k \in Z_2} \varepsilon_k/k = a/b$, then independent of the choice of the ε_k 's we may choose $b = \text{l.c.m. } Z_2$. From the definitions of $\psi(x)$ and $\vartheta(x)$ [2, pp. 340-341] we deduce that $\log b = \psi(N) - (\vartheta(N) - \vartheta(Q))$. Since $\psi(x) = \sum_{k=1}^x \vartheta(x^{1/k})$, one can show $\psi(x) - \vartheta(x) < 1.5x^{1/2}$ (see [4, Theorem 13]). Hence we see that $\log b \leq \vartheta(Q) + 1.5\sqrt{N}$. On the other hand

$$\frac{a}{b} \leq \sum_{i=1}^N \frac{1}{i} \leq \log N + \gamma + \frac{1}{N}$$

where $\gamma = 0.57 \dots$ is Euler's constant. Thus we see that the number of distinct possibilities for a is at most $b(\log N + \gamma + 1/N)$. It follows that

$$S_2(N) \leq (\log N + \gamma + 1/N) \exp(\vartheta(Q) + 1.5\sqrt{N}).$$

Whence

$$(2) \quad S_2^*(N) \leq \log(\log N + \gamma + 1/N) + \vartheta(Q) + 1.5\sqrt{N}.$$

Since $S^*(N) \leq S_1^*(N) + S_2^*(N)$ we can now estimate $S^*(N)$.

By the above estimates (1) and (2) for $S_1^*(N)$ and $S_2^*(N)$ we get

$$\begin{aligned} S^*(N) \leq \frac{N}{\log_2 N} \left\{ \log 2 \left((\log_2 N)^2 - \log_2 Q \log_2 N + \frac{\log_2 N}{(\log N)^2} + \frac{\log_2 N}{2(\log Q)^2} \right) \right. \\ \left. + \frac{\log(\log N + \gamma + 1/N) \cdot \log_2 N}{N} \right. \\ \left. + \frac{1.02 \log_2 N}{\log N} + \frac{1.5 \log_2 N}{\sqrt{N}} \right\} \end{aligned}$$

where we have used [4, Theorem 9] for the penultimate term. A straightforward calculation shows that for $\log_2 N \geq 1$ the term in the braces is decreasing when $N \geq 10^8$, and is less than 1.

Subcase B. $10^8 \geq N \geq e^e$. If $\log_2 N \leq 1/\log 2 = 1.4 \dots$, i.e., $N \leq 68.8 \dots$, then $2^N \leq \exp(N/\log_2 N)$ and the desired inequality holds.

For $N = 69, 70, 71, 72$, or 73 we note by direct calculation from the definition that $|Z_1| \leq 23 \leq N \cdot (23/69) = N/3$. Thus

$$S_1^*(N) \leq \frac{N \log 2}{3} \leq \frac{N \log 2}{\log_2 N} \cdot \left(\frac{1}{2}\right),$$

$$\begin{aligned} S_2^*(N) &\leq \log(\log(N + \gamma + 1/N) + \vartheta(Q) + 1.5\sqrt{N}) \\ &\leq \frac{N}{\log_2 N} \left\{ \frac{\log(\log(N + \gamma + 1/N)) \log_2 N}{N} + \frac{\log_2 N}{\log N} \frac{1.5 \log_2 N}{\sqrt{N}} \right\}. \end{aligned}$$

Since $S^*(N) \leq S_1^*(N) + S_2^*(N)$ we obtain

$$S^*(N) \leq \frac{N}{\log_2 N} \left\{ \frac{\log 2}{2} + \frac{\log(\log(N + 1)) \log_2 N}{N} + \frac{\log_2 N}{\log N} + \frac{1.5 \log_2 N}{\sqrt{N}} \right\}.$$

Since the term in braces is less than 1 for $69 \leq N < 74$, the inequality holds for $N < 74$.

For $74 \leq N \leq 10^8$ we use the estimates of [4, Theorems 18, 20, and 13] to obtain the desired result in a manner analogous to the case when $N \geq 10^8$. The difference in the cases $74 \leq N \leq 10^8$ and $N \geq 10^8$ are all consequences of the different estimates for $\sum 1/p$ and $\vartheta(x)$. The calculations are left to the reader.

Thus the first half of Lemma 8 is established.

The case $\log_4 N \geq 1$. In this case $N \geq 10^8$. From (1) and (2) we get

$$\begin{aligned} S^*(N) &\leq \frac{N \log_2 N}{\log N} \left\{ \log 2 \left(\log N - \frac{\log_2 Q \log N}{\log_2 N} \right. \right. \\ &\quad \left. \left. + \frac{1}{\log N \log_2 N} + \frac{\log N}{2 \log_2 N \log^2 Q} \right) \right. \\ &\quad \left. + \left(\frac{\log(\log N + 1) \log N}{N \log_2 N} + \frac{1.02}{\log_2 N} + \frac{1.5 \log N}{\sqrt{N} \log_2 N} \right) \right\}. \end{aligned}$$

Using the estimates

$$\log N - \frac{\log_2 Q \log N}{\log_2 N} \leq 1 + \frac{\log_2 N}{\log N}$$

in the above inequality yields

$$\begin{aligned} S^*(N) &\leq \frac{N \log_2 N}{\log N} \left\{ \log 2 \left(1 + \frac{\log_2 N}{\log N} + \frac{1}{\log N \log_2 N} + \frac{\log N}{2 \log_2 N \log^2 Q} \right) \right. \\ &\quad \left. + \left(\frac{\log(\log N + 1) \log N}{\log N \log_2 N} + \frac{\log^2 N}{N} + \frac{1.02}{\log_2 N} + \frac{1.5 \log N}{\sqrt{N} \log_2 N} \right) \right\}. \end{aligned}$$

An easy calculation shows that in the range under consideration, $\log_4 N \geq 1$, each term in the parentheses is decreasing. Trivial numerical estimates show that for $\log_4 N = 1$ the quantity in braces is less than 1.

Lemma 8 is proved.

LEMMA 9. Let $Q = N/\log N$ and $Q' = N/\log_2 N$. Suppose that $\log_6 N \geq 1$. Then

$$\sum_{Q < p \leq Q'} \frac{1}{p \log(N/p)} \leq \frac{\log_3 N}{\log N} \left(1 - \frac{\log_4 N}{2 \log_3 N} \right).$$

Proof. This is proved by using Lemma 6 almost exactly the same way it was used in the paragraphs following its proof, except that in this case $f(x)$ is simpler and slight adjustments must be made since we are deriving an upper bound.

The details are left to the reader.

THEOREM 3. For $r \geq 1$ and $\log_{2^r} N \geq 1$,

$$S(N) \leq \exp \left(\frac{N \log_r N}{\log^2 N \log_2 N} \prod_{j=1}^r \log_j N \right).$$

Proof. The values $r = 1, 2$ yield the statements of Lemma 8. We suppose the result is true for $r - 1 \geq 2$ and show that it holds for r .

We divide the integers less than N in a way similar to that in the proof of Theorem 2. Let $Q = N/\log N$ and $Q' = N/\log_2 N$. We define Z_1 and Z_2 by

$$Z_1 = \{k \leq N: \text{there exists } p, Q < p < N, p | k\}$$

and

$$Z_2 = \{k \leq N: k \notin Z_1\}.$$

Thus

$$\sum_{k=1}^N \frac{\varepsilon_k}{k} = \sum_{k \in Z_1} \frac{\varepsilon_k}{k} + \sum_{k \in Z_2} \frac{\varepsilon_k}{k}.$$

If $S_i(N)$ denotes the number of distinct values of the sums over Z_i as the ε_k 's take on all possible values with $\varepsilon_k = 0$ or 1, then $S(N) \leq S_1(N)S_2(N)$. We estimate each of $S_1(N)$ and $S_2(N)$ separately. Let $S_1^*(N) = \log S_1(N)$; then $S^*(N) \leq S_1^*(N) + S_2^*(N)$.

We estimate $S_2^*(N)$ first. For any choice of ε_k 's we may write

$$\sum_{k \in Z_2} \frac{\varepsilon_k}{k} = \frac{a}{b} \quad \text{where } a < \left(\sum_{i=1}^N \frac{1}{i} \right) b \text{ and } b = \text{l.c.m.}(Z_2).$$

As in the proofs of Lemma 8, we obtain from (2),

$$\begin{aligned} S_2^*(N) &\leq \log(\log N + 1) + \vartheta(Q) + 1.5\sqrt{N} \\ (4) \quad &\leq \log_2 N + 1/\log N + N/\log N + N/\log^2 N + 1.5\sqrt{N} \\ &\leq 2N/\log N \end{aligned}$$

where we have used [4, Theorem 4] and $1/(2 \log Q) < 1/\log N$ for the values of N under consideration.

We now turn to an estimation of $S_1(N)$. We rewrite the sum as follows

$$\sum_{k \in Z_1} \frac{\varepsilon_k}{k} = \sum_{Q < p < N} \frac{1}{p} \left(\sum_{k=1}^{N/p} \frac{\varepsilon_k}{k} \right)$$

where the ε_k 's on the internal sums (which properly should be $\varepsilon_{p,k}$) are independently taking on all possible combinations of values of 0 or 1. We see from this representation that

$$S_1^*(N) \leq \sum_{Q < p \leq N} S^*(N/p).$$

We break the sum in two parts as follows:

$$(5) \quad \Sigma_1 = \sum_{Q < p \leq Q'} S^*(N/p), \quad \Sigma_2 = \sum_{Q' < p \leq N} S^*(N/p).$$

Notice that for $Q < p \leq Q'$ we have $N/p \geq \log_2 N$ and thus

$$\log_{2(r-1)} N/p \geq \log_{2r} N \geq 1$$

so that the induction hypothesis for $r - 1$ is satisfied for N/p in the first sum. For the second sum we will use the estimates of Lemmas 7 and 8 which yield $S^*(x) \leq x \log 2$ and $S^*(x) \leq (x \log_2 x)/\log x$. We estimate Σ_2 first.

$$\Sigma_2 \leq \sum_{Q' < p \leq N/E} \frac{N \log_2 N/p}{p \log N/p} + \sum_{N/E < p \leq N} \frac{N}{p} \log 2$$

where E is chosen so that $\log_4 E = 1$. The first sum can be estimated by the use of Lemma 6 with

$$f(p) = \frac{\log_2(N/p)}{p \log(N/p)}.$$

After some calculation one gets

$$\sum_{Q' < p < N/E} f(p) \leq \frac{N \log_4^2 N}{\log N}.$$

Using the standard estimates [4, Theorem 5] for $\sum 1/p$ one obtains

$$\sum_{N/E < p \leq N} \frac{N}{p} \log 2 \leq \frac{N \log E}{\log N}.$$

We thus obtain

$$(6) \quad \Sigma_2 \leq \frac{N}{\log N} (\log_4^2 N + \log E).$$

We now estimate Σ_1 from (5), where we substitute for $S^*(N/p)$ the bound given by the induction hypothesis to obtain

$$(7) \quad \begin{aligned} \Sigma_1 &\leq \sum_{Q < p \leq Q'} \frac{N \log_{r-1}(N/p)}{p \log^2(N/p) \log_2(N/p)} \prod_{j=1}^{r-1} \log_j(N/p) \\ &< \frac{N \log_{r-1} N/Q}{\log N/Q \log_2 N/Q} \prod_{j=1}^{r-1} \log_j(N/Q) \sum_{Q < p \leq Q'} \frac{1}{p \log N/p}, \end{aligned}$$

where we have used the fact that

$$\frac{\log_{r-1} N/x}{\log N/x \log_2 N/x} \prod_{j=1}^{r-1} \log_j(N/x)$$

is decreasing in the interval $Q \leq x \leq Q'$ since the two terms in the denominator cancel into the numerator and the rest of the numerator is clearly decreasing in x . But $N/Q = \log N$ and $\sum 1/(p \log N/p)$ can be estimated by Lemma 9; thus

$$\Sigma_1 < \frac{N \log_r N}{\log_2 N \log_3 N} \left(\prod_{j=2}^r \log_j N \right) \frac{\log_3 N}{\log N} \left(1 - \frac{\log_4 N}{2 \log_3 N} \right).$$

The above can be rewritten as

$$(8) \quad \Sigma_1 < \frac{N \log_r N}{\log^2 N \log_2 N} \left(\prod_{j=1}^r \log_j N \right) \left(1 - \frac{\log_4 N}{2 \log_3 N} \right).$$

We combine (4), (6), and (8) to obtain

$$(9) \quad S^*(N) \leq \frac{N \log_r N}{\log N} \left(\prod_{j=3}^r \log_j N \right) \times \left(1 - \frac{\log_4 N}{2 \log_3 N} + \frac{\log_4^2 N + \log E}{\log_r N \prod_{j=3}^r \log_j N} + \frac{2}{\log_r N \prod_{j=3}^r \log_j N} \right).$$

It is not difficult to verify that the quantity in braces in (9) is less than 1; hence,

$$(10) \quad S^*(N) < \frac{N \log_r N}{\log N} \prod_{j=3}^r \log_j N.$$

But (10) is clearly equivalent to the inequality of Theorem 3, which is thus proven.

IV. A lower bound for $D(P)$

The proof is virtually the same as that for Theorem 2 of [1] except that we have a better bound for $S(N)$.

THEOREM 4. *If P is a prime then for P large enough that $\log_{2r} P \geq 1$*

$$D(P) \geq \frac{P \cdot \log P \cdot \log_2 P}{\log_{r+1} P \prod_{j=4}^{r+1} \log_j P}.$$

Proof. For each $a|P$, $1 \leq a < P$, write

$$\frac{a}{P} = \frac{1}{P} \left(\frac{1}{x_1} + \frac{1}{x_2} + \cdots + \frac{1}{x_{r_a}} \right) + \frac{1}{y_1} + \frac{1}{y_2} \cdots + \frac{1}{y_{s_a}}$$

where $x_i < x_{i+1}$, $(x_i, P) = (y_i, P) = 1$, and x_{i_n} is minimal for all expansions of a/P . Let $N = \max \{x_{i_n} : 1 \leq a < P\}$. Each value of a requires a different value of

$$\frac{1}{x_1} + \frac{1}{x_2} + \cdots + \frac{1}{x_{i_n}} = \sum_{k=1}^N \frac{\varepsilon_k}{k}$$

for some choice of ε_k 's. Thus N must be such that $S(N) \geq P$, the value $a = 0$ corresponding to the choice of all $\varepsilon_k = 0$. From Theorem 3 we see that for P large enough that $\log_{2r} P \geq 1$, N must be bigger than

$$\frac{\log P \cdot \log_2 P}{\log_{r+1} P \prod_{j=4}^{r+1} \log_j P},$$

since for that value $S^*(N) < \log P$. The desired inequality follows.

There are both heuristic and experimental reasons to suppose that the order of $D(N)/N$ is largest for $N = P$, a prime. This could be established if one could prove that for $(M, N) = 1$, $D(MN) \leq D(M) \cdot D(N)$, since we already know [1, Theorem 5] that $D(P^k) \leq 2D(P)P^{k-1}$. Exact estimates for $D(P)$ seem difficult since $D(P)/P$ is not monotone.

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