

On the divisibility properties of sequences of integers (I)

by

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Let $a_1 < a_2 < \dots$ be a sequence A of integers. Put $A(x) = \sum_{a_i \leq x} 1$.

The sequence is said to *have positive lower density* if

$$\lim_{x \rightarrow \infty} (A(x)/x) > 0,$$

it is said to *have positive upper logarithmic density* if

$$\overline{\lim}_{x \rightarrow \infty} \frac{1}{\log x} \sum_{a_i \leq x} \frac{1}{a_i} > 0.$$

The definition of upper density and lower logarithmic density is self-explanatory.

Besicovitch ([2]) was the first to construct a sequence of positive upper density no term of which divides any other. Behrend ([1]) and Erdős ([4]) on the other hand proved that in a sequence of positive lower density there are infinitely many couples satisfying $a_i | a_j$, Behrend in fact proved this if we only assume that the upper logarithmic density is positive.

Davenport and Erdős ([3]) proved that if A has positive upper logarithmic density there is an infinite subsequence $a_{i_j}, 1 \leq j < \infty$ satisfying $a_{i_j} | a_{i_{j+1}}$.

Put

$$f(x) = \sum_{\substack{a_i | a_j \\ a_j < x}} 1.$$

It is reasonable to conjecture that if A has positive density then

$$(1) \quad \lim_{x \rightarrow \infty} \frac{f(x)}{x} = \infty.$$

We have proved (1) and in fact obtained a fairly accurate determination of the speed with which $f(x)/x$ has to tend to infinity, this

strongly depends on the numerical value of the density of A . We will prove (1) in a subsequent paper.

Throughout this paper c_1, c_2, \dots will denote positive absolute constants, not necessarily the same at each occurrence, $\log_k x$ denotes the k -fold iterated logarithm. In the present paper we shall prove the following

THEOREM 1. *Assume that the sequence A has positive upper logarithmic density and put*

$$(2) \quad \overline{\lim} \frac{1}{\log x} \sum_{a_i < x} \frac{1}{a_i} = c_1.$$

Then there is a c_2 depending only on c_1 so that for infinitely many x

$$(3) \quad f(x) > x e^{c_2(\log_2 x)^{1/2} \log_3 x}.$$

On the other hand there is a sequence A satisfying (2) so that for all x

$$(4) \quad f(x) < x e^{c_3(\log_2 x)^{1/2} \log_3 x}.$$

First we prove (3). Our principal tool will be the following purely combinatorial

THEOREM 2. *Let \mathcal{S} be a set of n elements and let $B_1, \dots, B_z, z > c_4 2^n$ ($c_4 < 1$) be subsets of \mathcal{S} . Then if $n > n_0(c_4)$ one of the B 's contains at least $e^{c_5 n^{1/2} \log n}$ of the B 's, where c_5 depends only on c_4 .*

Before we prove Theorem 2 we show that apart from the value of c_5 it is best possible. To see this let the B 's be all subsets of \mathcal{S} having t elements where $\frac{1}{2}n + c_6 n^{1/2} > t > \frac{1}{2}n - c_6 n^{1/2}$. A simple computation shows that for suitable $c_6, z > c_4 2^n$ and every B contains fewer than $e^{c_7 n^{1/2} \log n}$ other B 's.

To prove Theorem 2 we first note the well known fact that for suitable c_8

$$(5) \quad \sum_1 \binom{n}{j} + \sum_2 \binom{n}{j} < \frac{c_4}{2} 2^n,$$

where in $\sum_1, j < \frac{1}{2}n - c_8 n^{1/2}$ and in $\sum_2, j > \frac{1}{2}n + c_8 n^{1/2}$. Because of (5) we can assume without loss of generality (replacing c_4 by $\frac{1}{2}c_4$) that $|B|$ denotes the number of elements of B

$$(6) \quad \frac{1}{2}n - c_8 n^{1/2} < |B_i| < \frac{1}{2}n + c_8 n^{1/2}.$$

Denote by $\mathcal{S}^{(j)}$ the family of these B 's which have precisely j elements (j satisfies (6)) and denote by $B_1^{(j)}, \dots, B_{g^{(j)}}^{(j)}$ the sets of $\mathcal{S}^{(j)}$. Clearly

$$(7) \quad \sum' g^{(j)} g(j) \leq \frac{c_4}{2} 2^n \leq \frac{z}{2},$$

where in \sum' the summation is extended over those j 's for which $g(j) \leq \frac{c_4}{2} \binom{n}{j}$. By (7) and $\binom{n}{j} < \frac{c_2 2^n}{\sqrt{n}}$ we can assume without loss of generality that either $g(j) = 0$ or $g(j) > \frac{1}{2} c_4$ and that

$$(8) \quad \sum g(j) > c_9 \sqrt{n}.$$

We obtain this by considering only the B 's which have j elements where $g(j) > \frac{1}{2} c_4$.

Put

$$s = \left[\frac{2}{c_4} \right] + 2.$$

From (8) we obtain by a simple argument that for a suitable c_{10} there is a sequence $j_1 < j_2 < \dots < j_s$ satisfying

$$(9) \quad g(j_r) > \frac{1}{2} c_4, \quad r = 1, \dots, s$$

and

$$(10) \quad j_{r+1} - j_r > c_{10} n^{1/2}, \quad r = 1, \dots, s-1.$$

From (10) we obtain by a simple computation that

$$(11) \quad \binom{j_r}{j_{r-1}} > e^{c_{11} n^{1/2} \log n}, \quad r = 1, \dots, s.$$

We are going to show that c_5 can be chosen as $\frac{1}{2} c_{11}$. In fact we shall show that if we consider only the set of $\mathcal{S}^{(j_r)}$, $r = 1, \dots, s$ and denote these sets by B'_1, \dots, B'_{r_1} then there is a B' which contains at least

$$(12) \quad e^{c_5 n^{1/2} \log n}, \quad c_5 = \frac{1}{2} c_{11}$$

B 's. Assume that (12) is false for sufficiently large n , we will arrive at a contradiction. Denote by $I^{(j_r)}$ the subsets of \mathcal{S} having j_r elements which contain at least $e^{c_5 n^{1/2} \log n}$ of the sets B . By our assumption the families $I^{(j_r)}$ and $\mathcal{S}^{(j_r)}$ are disjoint. Denote $I^{(j_r)} \cup \mathcal{S}^{(j_r)} = V^{(j_r)}$. Put

$$|I^{(j_r)}| = h(j_r), \quad |V^{(j_r)}| = \varphi(j_r).$$

By our assumption we have

$$(13) \quad \varphi(j_r) = h(j_r) + |\mathcal{S}^{(j_r)}| \geq h(j_r) + \frac{1}{2} c_4 \binom{n}{j_r}.$$

We will obtain our contradiction by showing that for a suitable r

$$(14) \quad \varphi(j_r) > \binom{n}{j_r}.$$

Now we estimate $\varphi(j_r)$ from below. First of all we evidently have

$$(15) \quad \varphi(j_1) = |\mathcal{S}^{(j_1)}| > \frac{1}{2}c_4 \binom{n}{j_1}.$$

Now we show that for every $r \leq s$ ($s = \left\lceil \frac{2}{c_4} \right\rceil + 2$)

$$(16) \quad \varphi(j_r) > (r + o(1)) \frac{1}{2}c_4 \binom{n}{j_r}.$$

To prove (16) we use induction with respect to r . By (15), (16) holds for $r = 1$. Assume that it holds for $r-1$, we will deduce it for r . To show this we will prove that if (16) holds for $r-1$ then

$$(17) \quad h(j_r) > (r-1 + o(1)) \binom{n}{j_r}.$$

By (13), (17) implies (16) for r and thus we only have to prove (17). Consider now all the subsets of \mathcal{S} having j_r elements which contain one of the sets of $V^{(j_{r-1})}$. We will estimate $h(j_r)$ from below by counting in two ways the number of times a subset of \mathcal{S} having j_r elements can contain a set of $V^{(j_{r-1})}$. First of all there are clearly $\varphi(j_{r-1}) \binom{n-j_{r-1}}{j_r-j_{r-1}}$ such relations, since to each of the $\varphi(j_{r-1})$ sets of $V^{(j_{r-1})}$ there are clearly $\binom{n-j_{r-1}}{j_r-j_{r-1}}$ subsets of \mathcal{S} having j_r elements which contain it. On the other hand the $h(j_r)$ sets of $I^{(j_r)}$ each contain at most $\binom{j_r}{j_{r-1}}$ sets of $V^{(j_{r-1})}$ (since they contain at most $\binom{j_r}{j_{r-1}}$ subsets having j_{r-1} elements). The other $\binom{n}{j_r} - h(j_r)$ subsets of \mathcal{S} having j_r elements contain fewer than $e^{c_5 n^{1/2} \log n}$ sets of $V^{(j_{r-1})}$. To see this observe that such a set can not contain a set of $I^{(j_{r-1})}$ since otherwise it would belong to $I^{(j_r)}$ and since it does not belong to $I^{(j_r)}$ it contains fewer than $e^{c_5 n^{1/2} \log n}$ sets of $\mathcal{S}^{(j_r)}$. Thus we evidently have

$$(18) \quad \varphi(j_{r-1}) \binom{n-j_{r-1}}{j_r-j_{r-1}} < h(j_r) \binom{j_r}{j_{r-1}} + \binom{n}{j_r} e^{c_5 n^{1/2} \log n}.$$

From (18) we obtain by a simple computation using (11) and $c_5 = \frac{1}{2}c_{11}$

$$(19) \quad \begin{aligned} h(j_r) &> \varphi(j_{r-1}) \binom{n-j_{r-1}}{j_r-j_{r-1}} \binom{j_r}{j_{r-1}}^{-1} - \binom{n}{j_r} e^{c_5 n^{1/2} \log n} \binom{j_r}{j_{r-1}}^{-1} \\ &\geq \varphi(j_{r-1}) \binom{n}{j_{r-1}}^{-1} \binom{n}{j_r} - \binom{n}{j_r} e^{-c_5 n^{1/2} \log n}. \end{aligned}$$

In (19) we use

$$\binom{n-j_{r-1}}{j_r-j_{r-1}} \binom{j_r}{j_{r-1}}^{-1} = \binom{n}{j_{r-1}}^{-1} \binom{n}{j_r}.$$

From (19) and the fact that (16) holds for $r-1$ we have

$$h(j_r) > (r-1 + o(1)) \binom{n}{j_r},$$

which proves (17), and hence (16) holds for all $r \leq s$.

But (16) implies that (14) holds for $r = s$. This contradiction proves Theorem 2.

By the same method we would prove the following

THEOREM 3. *Let \mathcal{S} be a set of n elements and let $B_1, \dots, B_z, z > c \frac{2^n}{\sqrt{n}} x$, where $x > 1, z \leq 2^n$ and c is a sufficiently large constant. Then if $n > n_0$ one of the B' contains at least $e^{c'x \log n}$ of the B 's.*

Theorem 3 clearly contains Theorem 2. The proof of Theorem 3 is similar but somewhat more complicated than that of Theorem 2. We suppress the proof of Theorem 3.

The proof of (3) is now a simple task. In fact we shall prove the following slightly stronger

THEOREM 1'. *Let $a_1 < \dots < a_l \leq N$ be a sequence of integers satisfying*

$$(20) \quad \sum_{i=1}^l \frac{1}{a_i} > c_{12} \log N.$$

Then there is a constant c_{13} depending only on c_{12} so that if $N > N_0(c_{11}, c_{12})$ then

$$(21) \quad \sum^+ \frac{1}{a_i} > \frac{1}{2} c_{12} \log N$$

where in (21) the summation is extended over the a 's, which have at least $\exp(c_{13}(\log_2 N)^{1/2} \log_3 N)$ divisors among the a 's.

It is easy to see that Theorem 1' implies Theorem 1. To see this observe that if (2) holds then (20) holds for infinitely many N . But if (21) holds a simple computation shows that to each N which satisfies (21) there is an $M = M(N) < N$ which tends to infinity with N and for which the number of $a_l < M$ which have at least $\exp(c_{13}(\log_2 N)^{1/2} \log_3 N)$ divisors among the a 's is greater than $\frac{1}{4} c_{12} M$. Thus M satisfies (3) and hence Theorem 1' implies (3).

Thus we only have to prove Theorem 1'. Assume that Theorem 1' is false. Then for arbitrarily large values of n there exists a sequence

$a_1 < \dots < a_l \leq N$ satisfying (20) which does not satisfy (21). Then there clearly exists a subsequence of the sequence $a_1 < \dots$, say $b_1 < \dots < b_r \leq N$ satisfying

$$(22) \quad \sum_{i=1}^r \frac{1}{b_i} > \frac{1}{2} c_{12} \log N$$

so that each b has fewer than $\exp(c_{13}(\log_2 N)^{1/2} \log_3 N)$ divisors among the b 's. We now show that this conclusion leads to a contradiction.

First we observe that by using

$$\sum_{k=1}^{\infty} \frac{1}{k^2} = \frac{\pi^2}{6} < 2$$

we obtain that there is a t so that there is a subsequence $b_{i_1} < \dots < b_{i_s}$ of the b 's each of which can be written in the form

$$b_{i_r} = t^2 q_r, \quad 1 \leq r \leq s$$

where the q_r are squarefree integers and where

$$(23) \quad \sum_{r=1}^s \frac{1}{q_r} > \frac{1}{4} c_{12} \log N.$$

(23) immediately follows from the fact that every integer can be written (uniquely) as the product of a square and a squarefree number.

$d(n)$ (as usual) will denote the number of divisors of n . $d^+(n)$ denotes the number of q 's which divide n . By our assumption we have for all r ($r = 1, \dots, s$)

$$(24) \quad d^+(q_r) < \exp(c_{13}(\log_2 N)^{1/2} \log_3 N).$$

From (23) we have for $N > N_0$

$$(25) \quad \sum_{m=1}^N d^+(m) = \sum_{r=1}^s \left[\frac{N}{q_r} \right] \geq N \sum_{r=1}^s \frac{1}{q_r} - N > \frac{1}{5} c_{12} N \log N.$$

Denote by $\nu(m)$ the number of distinct prime factors of m . Since the q 's are squarefree we have $d^+(n) \leq 2^{\nu(n)}$.

Thus from (25) we obtain (the dash indicates that the summation is extended over the $n \leq N$ for which $\nu(n) > \log_2 N$)

$$(26) \quad \sum_{m=1}^N d^+(m) > \frac{1}{5} c_{12} N \log N - N 2^{\log_2 N} > \frac{1}{10} c_{12} N \log N.$$

On the other hand we evidently have

$$\sum_{m=1}^N d(m) = \sum_{m=1}^N \left[\frac{N}{m} \right] < 2N \log N.$$

Thus by (26) there is an m satisfying $v(m) > \log_2 N$ for which

$$(27) \quad d^+(m) > \frac{c_{12}}{20} d(m) \geq \frac{1}{20} c_{12} 2^{v(m)}.$$

The last equality of (27) follows from the fact that since the q 's are square-free we can assume that m is squarefree.

Now we can apply Theorem 2. The set \mathcal{S} is the set of prime divisors of m , $v(m) = n$. The B 's are the q 's which divide m , $c_{12}/20 = c_4$. We thus obtain by Theorem 2 that there is a q/m for which

$$\bar{d}^+(q) > \exp(c_5(\log_2 N)^{1/2} \log_3 N)$$

which contradicts (24) if c_{13} is sufficiently small.

This completes the proof of Theorem 1' and hence (3) is proved. It is clear from the above proof that (21) would remain true with $1-\varepsilon$ instead of $\frac{1}{2}$.

To complete the proof of Theorem 1 we now have to show (4). (We do not give the proof in full detail.) In fact we shall prove the following stronger

THEOREM 4. *There is an infinite sequence A of positive density for which for all x*

$$(28) \quad f(x) < x \exp(c_{14}(\log_2 x)^{1/2} \log_3 x).$$

Our principal tool for the proof of Theorem 4 will be the following result from probabilistic number theory:

THEOREM 5. *Let n be squarefree. Let $n = \prod_k p_k^{(n)}$, $p_1^{(n)} < \dots < p_{v(n)}^{(n)}$, be the decomposition of n into primes. Then for every $c_{15} > 0$ there is a $k_0 = k_0(c_{15})$ so that the density of integers n which satisfy for all $k_0 < k \leq v(n)$*

$$(29) \quad e^{e^k - c_{15}(\log_2 n)^{1/2}} < p_k < e^{e^k + c_{15}(\log_2 n)^{1/2}}$$

is positive.

Theorem 5 can be proved by the methods of probabilistic number theory ([5], [6]). We do not give here the proof of Theorem 5.

Now we show that the sequence of integers which satisfy (29) for all $k > k_0(c_{15})$ also satisfy (28) and if this is accomplished Theorem 4 and therefore (4) is proved. Thus the proof of Theorem 1 will be complete.

Let $a_1 < \dots < a_l \leq x$ be the sequence of integers satisfying (29). From (29) we obtain by a simple computation that for every r , $1 \leq r \leq l$

$$(30) \quad \log_2 a_r - 2c_{14}(\log_2 a_r)^{1/2} < v(a_r) < \log_2 a_r + 2c_{14}(\log_2 a_r)^{1/2}.$$

Denote as before by $d^+(a_r)$ the number of a 's dividing a_r . To prove (28) it will suffice to show that for every r

$$(31) \quad d^+(a_r) < \exp(c_{14}(\log_2 x)^{1/2} \log_3 x).$$

Denote by $p_1 < \dots < p_{r(a_r)}$ the prime factors of a_r . Assume $a_t | a_r$. If $v(a_t) \leq k_0$ then by (30) there are clearly fewer than $v(a_r)^{k_0+1} \leq (\log_2 x)^{k_0+2}$ choices for a_t , thus these can be ignored. If $v(a_t) > k_0$, let p_s be the greatest prime factor of a_t . Since a_t and a_r both satisfy (29) and (30) a simple computation shows that

$$(32) \quad s - 3c_{14}(\log_2 a_r)^{1/2} \leq v(a_t) \leq s.$$

Thus by an easy argument and simple computation

$$\begin{aligned} d^+(a_r) &\leq (\log_2 x)^{k_0+2} + \sum_{s=k_0+1}^{v(a_r)} \sum_{s-3c_{14}(\log_2 a_r)^{1/2}}^s \binom{s}{u} \\ &< (\log_2 x)^{k_0+2} + v(a_r) (v(a_r))^{4c_{14}(\log_2 a_r)^{1/2}} \\ &< v(a_r)^{5c_{14}(\log_2 a_r)^{1/2}} < \exp(c_{16}(\log_2 x)^{1/2} \log_3 x). \end{aligned}$$

Thus (31) is proved (with $c_{16} = c_{14}$).

References

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