

# ON THE STRENGTH OF CONNECTEDNESS OF A RANDOM GRAPH

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Let  $G$  be a non-oriented graph without parallel edges and without slings, with vertices  $V_1, V_2, \dots, V_n$ . Let us denote by  $d(V_k)$  the *valency* (or degree) of a point  $V_k$  in  $G$ , i. e. the number of edges starting from  $V_k$ . Let us put

$$(1) \quad c(G) = \min_{1 \leq k \leq n} d(V_k).$$

If  $G$  is an arbitrary non-complete graph, let  $c_p(G)$  denote the least number  $k$  such that by deleting  $k$  appropriately chosen vertices from  $G$  (i. e. deleting the  $k$  points in question and all edges starting from these points) the resulting graph is not connected. If  $G$  is a complete graph of order  $n$ , we put  $c_p(G) = n - 1$ . Let  $c_e(G)$  denote the least number  $l$  such that by deleting  $l$  appropriately chosen edges from  $G$  the resulting graph is not connected. We may measure the strength of connectedness of  $G$  by any of the numbers  $c_p(G)$ ,  $c_e(G)$  and in a certain sense (if  $G$  is known to be connected) also by  $c(G)$ . Evidently one has

$$(2) \quad c(G) \cong c_e(G) \cong c_p(G).$$

It is known further that any two points of  $G$  are connected by at least  $c_p(G)$  paths having no point in common, except the two endpoints (theorem of MENGER—WHITNEY, see [1] and [2]) and by at least  $c_e(G)$  paths having no edge in common (theorem of FORD and FULKERSON, see [3]).

We shall denote by  $\nu_r(G)$  the number of vertices of  $G$  which have the valency  $r$  ( $r = 0, 1, 2, \dots$ ).

As in two previous papers ([4], [5]) we consider the random graph  $\Gamma_{n, N}$  defined as follows: Let there be given  $n$  labelled points  $V_1, V_2, \dots, V_n$ . Let us choose at random  $N$  edges among the  $\binom{n}{2}$  possible edges connecting these  $n$  points, so that each of the  $\binom{\binom{n}{2}}{N}$  possible choices of these edges should be equiprobable. We denote by  $\Gamma_{n, N}$  the random graph thus obtained. We shall denote by  $\mathbf{P}(\cdot)$  the probability of the event in the brackets.

The aim of this note is to investigate the strength of connectedness of the random graph  $\Gamma_{n,N}$  when  $n$  and  $N$  both tend to  $+\infty$ ,  $N=N(n)$  being a function of  $n$ . As it has been shown in [4], the following theorem holds:

**THEOREM 1.** *If we have  $N(n) = \frac{1}{2}n \log n + \alpha n + o(n)$  where  $\alpha$  is a real constant, then the probability of  $\Gamma_{n,N(n)}$  being connected tends to  $\exp(-e^{-2\alpha})$  for  $n \rightarrow +\infty$ .*

In this paper we shall prove the following theorem:

**THEOREM 2.** *If we have  $N(n) = \frac{1}{2}n \log n + \frac{r}{2}n \log \log n + \alpha n + o(n)$  where  $\alpha$  is a real constant and  $r$  a non-negative integer, then*

$$(3) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(c_p(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2\alpha}}{r!}\right),$$

further

$$(4) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(c_e(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2\alpha}}{r!}\right)$$

and

$$(5) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(c(\Gamma_{n,N(n)}) = r) = 1 - \exp\left(-\frac{e^{-2\alpha}}{r!}\right).$$

**REMARK.** Clearly Theorem 2 can be considered as a generalization of Theorem 1. As a matter of fact, any of the statements  $c_p(G) = 0$  or  $c_e(G) = 0$  is equivalent to  $G$  not being connected and thus for  $r = 0$  (3) and (4) reduce to the statement of Theorem 1. It has been shown further in [4] that if  $N(n) = \frac{n}{2} \log n + \alpha n + o(n)$  and  $\Gamma_{n,N(n)}$  is not connected, then it consists almost surely of a connected component and of a few isolated points. Therefore (5) is for  $r = 0$  also equivalent to the statement of Theorem 1. Thus in proving Theorem 2 we may restrict ourselves to the case  $r \geq 1$ .

The statement (5) of Theorem 2 gives information about the *minimal* valency of points of  $\Gamma_{n,N}$ . In a forthcoming note we shall deal with the same question for larger ranges of  $N$  (when  $c(\Gamma_{n,N})$  tends to infinity with  $n$ ), further with the related question about the *maximal* valency of points of  $\Gamma_{n,N}$ .

We shall prove further the following

**THEOREM 3.** *If we have  $N(n) = \frac{1}{2}n \log n + \frac{r}{2}n \log \log n + \alpha n + o(n)$  where  $\alpha$  is a real constant and  $r$  a non-negative integer, then we have*

$$(6) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(v_r(\Gamma_{n,N(n)}) = k) = \frac{\lambda^k e^{-\lambda}}{k!} \quad \text{for } k = 0, 1, \dots$$

where  $\lambda = \frac{e^{-2\alpha}}{r!}$ ; in other words, the distribution of  $\nu_r(\Gamma_{n, N(n)})$  tends to a Poisson distribution.

PROOF OF THEOREMS 2 AND 3. Let  $r \geq 1$  be an integer and  $-\infty < \alpha < +\infty$ . Let us suppose that

$$(7) \quad N(n) = \frac{1}{2} n \log n + \frac{r}{2} n \log \log n + \alpha n + o(n).$$

Let  $\Gamma_{n, N}$  be a random graph with the  $n$  vertices  $V_1, V_2, \dots, V_n$  and having  $N$  edges. Let  $P_k(n, N, r)$  denote the probability that by removing  $r$  suitably chosen points from  $\Gamma_{n, N}$  there remain two disjoint graphs, consisting of  $k$  and  $n-k-r$  points, respectively. We may suppose  $k < \left\lfloor \frac{n-r}{2} \right\rfloor$ . First we have clearly

$$P_k(n, N, r) \leq \binom{n}{r} \binom{n-r}{k} \frac{\binom{\binom{n}{2} - k(n-k-r)}{N}}{\binom{\binom{n}{2}}{N}}.$$

It follows by some obvious estimations that

$$(8) \quad \sum_{\substack{(r+3) \frac{\log n}{\log \log n} < k \leq \left\lfloor \frac{n-r}{2} \right\rfloor}} P_k(n, N(n), r) = O\left(\frac{1}{n}\right).$$

Now we consider the case  $k \leq (r+3) \frac{\log n}{\log \log n}$ . Let  $P_k^*(n, N, r)$  denote the probability that by removing  $r$  suitably chosen points (the set of which will be denoted by  $\mathcal{A}$ )  $\Gamma_{n, N}$  can be split into two disjoint subgraphs  $\Gamma'$  and  $\Gamma''$  consisting of  $k$  and  $n-k-r$  points, respectively, but that  $\Gamma_{n, N}$  can not be made disconnected by removing only  $r-1$  points. If  $\Gamma_{n, N}$  has these properties and if  $s$  denotes the number of edges of  $\Gamma_{n, N}$  connecting a point of  $\mathcal{A}$  with a point of  $\Gamma'$ , then we have clearly  $s \geq r$ . Otherwise, by definition,  $s \leq rk$ . Thus we have

$$(9) \quad P_k^*(n, N, r) \leq \sum_{s=r}^{rk} \binom{n}{r} \binom{n-r}{k} \binom{rk}{s} \frac{\binom{\binom{n}{2} - k(n-k)}{N-s}}{\binom{\binom{n}{2}}{N}}.$$

It follows that

$$(10) \quad \sum_{k=2}^{\left[ \frac{(r+3) \log n}{\log \log n} \right]} P_k^*(n, N(n), r) = O\left(\frac{1}{\log n}\right).$$

From (8) and (10) it follows that for  $n \rightarrow +\infty$

$$(11) \quad \mathbf{P}(c_p(\Gamma_{n, N(n)}) = r) \sim \mathbf{P}(c(\Gamma_{n, N(n)}) = r).$$

As a matter of fact, (8) and (10) imply that if by removing  $r$  suitably chosen points (but not by removing less than  $r$  points)  $\Gamma_{n, N(n)}$  can be split into two disjoint subgraphs  $\Gamma'$  and  $\Gamma''$  consisting of  $k$  and  $n-k-r$  points, respectively,

where  $k \leq \left\lfloor \frac{n-r}{2} \right\rfloor$ , then only the case  $k=1$  has to be considered, the probability of  $k > 1$  being negligibly small. It remains to prove (5). This can be done as follows. First we prove that

$$(12) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(c(\Gamma_{n, N(n)}) \leq r-1) = 0.$$

For  $r=1$  this follows already from Theorem 1. Thus we may suppose here  $r \geq 2$ . We have

$$\mathbf{P}(c(\Gamma_{n, N}) \leq r-1) \leq \sum_{h=1}^{r-1} n \binom{n-1}{h} \frac{\binom{n}{2} - (n-1)}{N-h}, \frac{\binom{n}{2}}{N},$$

and thus

$$(13) \quad \mathbf{P}(c(\Gamma_{n, N(n)}) \leq r-1) = O\left(\frac{1}{\log n}\right)$$

which proves (12).

Now let  $\nu_r(\Gamma_{n, N})$  denote the number of vertices of  $\Gamma_{n, N}$  which have the valency  $r$ . Then we have clearly by (12)

$$(14) \quad \mathbf{P}(c(\Gamma_{n, N(n)}) = r) \sim \mathbf{P}(\nu_r(\Gamma_{n, N(n)}) \neq 0).$$

Now evidently

$$(15) \quad \mathbf{P}(\nu_r(\Gamma_{n, N(n)}) \neq 0) = \sum_{j=1}^n (-1)^{j-1} S_j$$

where

$$(16) \quad S_j = \sum_{1 \leq k_1 < k_2 < \dots < k_j \leq n} \dots \sum \mathbf{P}(d(V_{k_1}) = r, d(V_{k_2}) = r, \dots, d(V_{k_j}) = r).$$

Evidently, if we stop after taking an even or odd number of terms of the

sum on the right-hand side of (15), we obtain a quantity which is greater or smaller, respectively, than the left-hand side of (15). Now clearly

$$\mathbf{P}(d(V_i) = r) = \binom{n-1}{r} \frac{\binom{n}{2} - (n-1)}{N(n) - r} \sim \frac{e^{-2\alpha}}{nr!},$$

and thus

$$(17) \quad \lim_{n \rightarrow +\infty} S_1 = \frac{e^{-2\alpha}}{r!}.$$

Now let us consider  $\mathbf{P}(d(V_{k_1}) = r, d(V_{k_2}) = r)$  where  $k_1 \neq k_2$ . If both  $V_{k_1}$  and  $V_{k_2}$  have valency  $r$ , three cases have to be considered: a) either  $V_{k_1}$  and  $V_{k_2}$  are not connected, and there is no point which is connected with both  $V_{k_1}$  and  $V_{k_2}$ ; b) or  $V_{k_1}$  and  $V_{k_2}$  are not connected, but there is a point connected with both; c)  $V_{k_1}$  and  $V_{k_2}$  are connected. We denote the probabilities of the corresponding subcases by  $\mathbf{P}_a(d(V_{k_1}) = r, d(V_{k_2}) = r)$ ,  $\mathbf{P}_b(d(V_{k_1}) = r, d(V_{k_2}) = r)$  and  $\mathbf{P}_c(d(V_{k_1}) = r, d(V_{k_2}) = r)$ , respectively. We evidently have

$$\mathbf{P}_a(d(V_{k_1}) = r, d(V_{k_2}) = r) = \frac{(n-2)!}{r!(n-2r-2)!} \frac{\binom{n}{2} - (2n-3)}{N(n) - 2r} \sim \left(\frac{e^{-2\alpha}}{n \cdot r!}\right)^2,$$

and thus

$$(18) \quad \sum_{1 \leq k_1 < k_2 \leq n} \mathbf{P}_a(d(V_{k_1}) = r, d(V_{k_2}) = r) \sim \frac{1}{2} \left(\frac{e^{-2\alpha}}{r!}\right)^2.$$

On the other hand (denoting by  $l$  the number of points which are connected with both  $V_{k_1}$  and  $V_{k_2}$ ), we have

$$(19) \quad \begin{aligned} & \mathbf{P}_b(d(V_{k_1}) = r, d(V_{k_2}) = r) = \\ & = \sum_{l=1}^r \frac{(n-2)!}{l!(r-l)!(n-2r+l-2)!} \frac{\binom{n}{2} - (2n-3)}{N(n) - 2r} = O\left(\frac{1}{n^3}\right). \end{aligned}$$

Similarly one has

$$(20) \quad \begin{aligned} & \mathbf{P}_c(d(V_{k_1})=r, d(V_{k_2})=r) = \\ & = \sum_{l=0}^{r-1} \frac{(n-2)!}{l!(r-l-1)!(n-2r+l)!} \frac{\binom{n}{2} - (2n-3)}{N(n)-2r} \frac{1}{\binom{n}{2}} = O\left(\frac{1}{n^4}\right). \end{aligned}$$

Thus we obtain

$$\lim_{n \rightarrow +\infty} S_2 = \frac{1}{2} \left( \frac{e^{-2\alpha}}{r!} \right)^2.$$

The cases  $j > 2$  can be dealt with similarly. Thus we obtain

$$(21) \quad \lim_{n \rightarrow +\infty} S_j = \frac{1}{j!} \left( \frac{e^{-2\alpha}}{r!} \right)^j \quad (j = 1, 2, 3, 4, \dots).$$

It follows from (16) and (21) that

$$(22) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(\nu_r(I_{n, N(n)}) \neq 0) = 1 - \exp\left(-\frac{e^{-2\alpha}}{r!}\right).$$

In view of (2), (11) and (14) Theorem 2 follows.

To prove Theorem 3 it is sufficient to remark that by the well-known formula of CH. JORDAN

$$(23) \quad \mathbf{P}(\nu_r(I_{n, N(n)}) = k) = \sum_{j=0}^{n-k} (-1)^j \binom{j+k}{j} S_{j+k},$$

and thus by (21), putting  $\lambda = \frac{e^{-2\alpha}}{r!}$ , we obtain for  $k = 0, 1, \dots$

$$(24) \quad \lim_{n \rightarrow +\infty} \mathbf{P}(\nu_r(I_{n, N(n)}) = k) = \frac{\lambda^k}{k!} \sum_{j=0}^{\infty} \frac{(-1)^j \lambda^j}{j!} = \frac{\lambda^k e^{-\lambda}}{k!}.$$

Thus Theorem 3 is proved.

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### References

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