

# List colorings with distinct list sizes, the case of complete bipartite graphs

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## Abstract

A graph  $G$  is *f-choosable* if for every collection of lists with list sizes specified by  $f$  there is a proper coloring using colors from the lists. The sum choice number,  $\chi_{sc}(G)$ , is the minimum of  $\sum f(v)$ , over all  $f$  such that  $G$  is *f-choosable*. In this paper we show that  $\chi_{sc}(G)/|V(G)|$  can be bounded while the minimum degree  $\delta_{\min}(G) \rightarrow \infty$ . (This is not true for the list chromatic number,  $\chi_{\ell}(G)$ ). Our main tool is to give tight estimates for the sum choice number for the complete bipartite graphs  $K_{a,q}$ .

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# 1 Average list sizes and planar graphs

Given a graph  $G$  and a list of colors  $L(v)$  for each vertex  $v \in V(G)$  we say that  $G$  is  $L$ -choosable if it is possible to choose  $\ell(v) \in L(v)$  for all  $v$  so that  $\ell$  is a proper coloring of  $G$ . The *choice number* (or *list chromatic number*)  $\chi_\ell$  is the minimum  $t$  such that every assignment  $L$  with  $|L(v)| \geq t$  for all  $v \in V$  the graph is  $L$ -choosable. It is well-known (Grötzsch's theorem) that

$$\chi_\ell(P) \leq 5 \tag{1}$$

for every planar graph  $P$ , and this is the best possible. Erdős, Rubin and Taylor (see, e.g., [1]) showed for the complete bipartite graph that

$$\chi_\ell(K_{q,q}) = \Theta(\log q). \tag{2}$$

However, if we allow distinct list sizes, then the average size can be smaller. For example, Thomassen's beautiful proof [6] for (1) gives that if  $P$  is an  $n$ -vertex planar graph,  $v_1, \dots, v_t$  are its outside vertices and the list sizes are

$$|L(v)| = \begin{cases} 1 & \text{for } v = v_1, \\ 2 & \text{for } v = v_2, \\ 3 & \text{for } v = v_3, \dots, v_t, \\ 5 & \text{otherwise,} \end{cases} \tag{3}$$

then  $P$  is  $L$ -choosable.

Consider a function  $f : V(G) \rightarrow \mathbb{N}$ . An  $f$ -assignment is an assignment of lists  $L(v)$  to the vertices  $v \in V(G)$  such that  $|L(v)| = f(v)$  for all  $v$ . The function  $f$  itself is *choosable* if  $G$  is  $L$ -choosable for all  $f$ -assignments  $L$ . We define the *sum choice number* of  $G$ , denoted  $\chi_{sc}(G)$ , to be the least  $k$  for which there exists a choosable  $f$  with  $\sum_{v \in V(G)} f(v) = k$ . Thomassen's theorem implies that  $\chi_{sc}(P) \leq 5n - 9$ .

In fact, more is true. It is easy to show (see, e.g., [5]) that for every graph

$$\chi_{sc}(G) \leq |V(G)| + |E(G)| \tag{4}$$

holds. Hence  $\chi_{sc}(P) \leq 4n - 6$ . Our first result is a slight improvement.

**Theorem 1.1** *Let  $P$  be an  $n$ -vertex planar graph. There exists an  $f : V(P) \rightarrow \mathbb{N}$  such that  $\sum f(v) = 4n - 6$ ,  $\max f(v) \leq 6$ , and  $P$  is  $f$ -choosable.  $\square$*

## 2 Graphs with small average list sizes, $K_{a,q}$

Sum choice numbers were introduced by Isaak in [4] who proved that if  $G$  is the line-graph of  $K_{2,q}$  then  $\chi_{sc}(G) = q^2 + \lceil 5q/3 \rceil$ . Various classes of graphs were investigated by Isaak in [5], by Berliner, Bostelmann, Brualdi, and Deaett [2] and by Heinold [3]. Two results are of particular interest to us.

**Theorem 2.1** (Berliner et al. [2])  $\chi_{sc}(K_{2,q}) = 2q + 1 + \lfloor \sqrt{4q+1} \rfloor$ .

**Theorem 2.2** (Heinold [3])  $\chi_{sc}(K_{3,q}) = 2q + 1 + \lfloor \sqrt{12q+4} \rfloor$ .

Our main result extends Theorems 2.1 and 2.2 to arbitrary  $a$ .

**Theorem 2.3** *There exist constants  $c_1$  and  $c_2$  such that for all  $a \geq 4$  and  $q \geq 50a^2 \log a$*

$$2q + c_1 a \sqrt{q \log a} \leq \chi_{sc}(K_{a,q}) \leq 2q + c_2 a \sqrt{q \log a}.$$

This implies that we can find a choosable  $f$  such that the average list size does not necessarily grow with the average degree. Indeed, with  $q$  tending to infinity, average degree of  $K_{a,q}$  approaches  $2a$ . We obtain

$$\lim_{a \rightarrow \infty, q \gg a^2 \log a} \frac{|E(K_{a,q})|}{a+q} = \infty, \quad \lim_{a \rightarrow \infty, q \gg a^2 \log a} \frac{\chi_{sc}(K_{a,q})}{a+q} = 2. \quad (5)$$

## 3 List chromatic number and average degree

Alon has shown in [1] that  $\chi_l$  depends heavily on the average degree.

**Theorem 3.1** (Alon, [1]) *For some constant  $c$ , every graph  $G$  with average degree  $d$  has  $\chi_l(G) \geq c \frac{\log d}{\log \log d}$ .*

One of the most interesting corollaries of our Theorem 2.3 is (5), that if different list sizes are allowed, the conclusion of Theorem 3.1 is no longer true. Sum-choice depends more on the structure of the graph than the list chromatic number.

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Throughout this paper,  $\log$  is the natural logarithm and the two partite sets of the complete bipartite graph  $K_{a,q}$  is denoted by  $A$  and  $Q$ , with  $|A| = a$  and  $|Q| = q$ .

## 4 Upper bound, there are choosable short lists

**Theorem 4.1** *Suppose that  $a, q \in \mathbb{N}$  with  $q > a > 3$ . Then*

$$\chi_{sc}(K_{a,q}) \leq 2q + a \lceil \sqrt{32q(1 + \log a)} \rceil.$$

**Proof.** To prove the upper bound, we present a function  $f$  with  $\sum_{v \in A \cup Q} f(v) \geq 2q + a \sqrt{32q(1 + \log a)}$  such that every  $f$ -assignment is choosable.

Define  $f$  as

$$f(v) = \begin{cases} r & \text{for } v \in A; \\ 2 & \text{for } v \in Q \end{cases}$$

where  $r$  will be defined later in (6). Let  $L$  be an arbitrary  $f$ -assignment, i.e.  $|L(v)| = f(v)$  for all  $v$ .

Consider  $S := \cup_{v \in A \cup Q} L(v)$ . The assignment  $L$  yields a multihypergraph and a multigraph on the vertex set  $S$  and with edge sets  $\mathcal{L}_A := \{L(u) : u \in A\}$  and  $\mathcal{L}_Q := \{L(v) : v \in Q\}$ , respectively. Choosability of  $L$  means that one can find a set  $T \subset S$  meeting all hyperedges of  $\mathcal{L}_A$  such that  $S \setminus T$  meets all edges of  $\mathcal{L}_Q$ , so  $T$  is an independent set in the graph  $\mathcal{L}_Q$ . Then the choice function  $\ell$  can be defined as

$$\ell(u) \in L(u) \cap T, \text{ for } u \in A$$

and

$$\ell(v) \in L(v) \cap (S \setminus T), \text{ for } v \in Q.$$

We are going to construct such a  $T$  by a 2-step random process.

Let us pick, randomly and independently, each element of  $S$  with probability  $p$ . Let  $B$  be the random set of all elements picked. Define a random variable  $X_u$  for each  $u \in A$  as  $X_u = |L(u) \cap B|$ , and the random variable  $Y$  by  $Y := |\{v \in Q : L(v) \subseteq B\}|$ , so  $Y$  is the number of edges of  $\mathcal{L}_Q$  spanned by  $B$ . Remove an element  $\ell(v) \in L(v)$  for each edge of  $\mathcal{L}_Q$  spanned by  $B$ , the remaining set  $T \subset B$  is certainly independent in  $\mathcal{L}_Q$ . If  $Y < X_u$  for each  $u \in A$ , then  $T$  meets all  $L(u) \in \mathcal{L}_A$  and we are done.

One needs a careful definition

$$p := \sqrt{\frac{2(1 + \log a)}{q}} \quad \text{and} \quad r \geq 4pq = \sqrt{32(1 + \log a)}q. \quad (6)$$

Standard probabilistic arguments (Chernoff inequality) complete the proof.  $\square$

## 5 Lower bound, much shorter lists are not choosable

To prove that  $\chi_{sc}(G) \geq k$  for a particular  $k$ , we need to show that for every  $f$  with  $\sum_{v \in G} f(v) = k$ , there exists a non-choosable  $f$ -assignment. Here we only show how to construct a non-choosable assignment for a very special  $f$ .

**Lemma 5.1** *Let  $t \geq 2$  and  $l \geq 1$ . For  $a = \binom{2t}{t}$  and  $q = t\ell^2$ , there exists a non-choosable assignment  $L$  with  $L(v) = 2$  for  $v \in Q$  and  $L(v) = t\ell$  for  $v \in A$ .*

Note that with this choice of  $a$  and  $q$ , we have  $|L(v)| \geq \sqrt{\frac{q \log_2 a}{2}}$  for  $v \in A$ .

**Proof.** Let us define the vertex set of a hypergraph  $\mathcal{H}$  as  $V(\mathcal{H}) = \cup_{i=1}^{2t} A_i$  where the  $A_i$ 's are disjoint  $\ell$ -sets. The edges of  $\mathcal{H}$  are of the form  $\cup_{i \in I} A_i$  for all subsets  $I \subseteq \{1, \dots, 2t\}$  of size  $t$ . Define a graph  $G$  on the vertex set  $\cup_{i=1}^{2t} A_i$ . Let  $\{x, y\}$  be an edge if and only if  $x \in A_{2i-1}$  and  $y \in A_{2i}$ .

Define the lists of the vertices  $v \in A$  to be the sets in  $E(\mathcal{H})$  and the two element sets in  $E(G)$  to be the lists of the vertices  $v \in Q$ .  $\square$

This argument can be extended to every  $f$  if  $\sum f(v)$  is sufficiently small.

**Theorem 5.2** *If  $a \geq 3$  and  $q > 50a^2 \log a$ , then*

$$\chi_{sc}(K_{a,q}) \geq 2q + 0.068a\sqrt{q \log a}. \quad \square$$

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