

# ON THE PRAGUE DIMENSION OF KNESER GRAPHS

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ABSTRACT. In this note we point out another connection between the Prague dimension of graphs and the dimension theory of partially ordered sets by giving a very short proof of a theorem of Poljak, Pultr and Rödl [10]. We show that the dimension of the Kneser graph is bounded as  $\dim_P(K(n, k)) < C_k \log \log n$ , where  $C_k$  is depending only on  $k$ .

## 1. DIMENSION OF GRAPHS

The Kneser graph  $K(n, k)$  is the graph whose vertices are the  $k$ -subsets of the  $n$ -element set  $[n] := \{1, 2, \dots, n\}$ , with vertices are adjacent when the corresponding  $k$ -sets are disjoint.

The *product* of the graphs  $(V_1, E_1)$  and  $(V_2, E_2)$  is a graph with vertex set  $V_1 \times V_2$ ; two vertices  $(v_1, v_2)$  and  $(w_1, w_2)$  are adjacent in the product graph if  $(v_1, w_1)$  is adjacent in  $G_1$  and  $(v_2, w_2)$  is adjacent in  $G_2$ . In particular,  $v_i$  and  $w_i$  must be distinct. The *Prague dimension* (or product dimension) of the graph  $G$ ,  $\dim_P(G)$ , is the minimum number  $d$  such that  $G$  is an induced subgraph of the product of  $d$  complete graphs. In other words, it is the minimum  $d$  such that the vertices  $x$  of  $G$  can be represented by vectors  $\mathbf{v}(x) = (v_1(x), \dots, v_d(x))$  such that  $(x, y)$  forms an edge if and only if  $v_i(x) \neq v_i(y)$  for all  $1 \leq i \leq d$ . Again, another form, it is the minimum number of good colorings of the vertices of  $G$ ,  $\varphi_1, \dots, \varphi_d$ , (not necessarily with minimum number of colors), such that for every non-edge  $(a, b)$  one has at least one  $i$  with  $\varphi_i(a) = \varphi_i(b)$ .

The Prague dimension was introduced and investigated in a series of papers by Nešetřil, Pultr [9], and other Czech mathematicians. Poljak, Pultr and Rödl [10] proved that

$$(1) \quad \log_2 \log_2 (n/(k-1)) \leq \dim_P(K(n, k)) \leq C_k \lceil \log_2 \lceil \log_2 n \rceil \rceil ,$$

with  $C_k \leq (k-1)k^2$ . Later (for  $n$  sufficiently large) they [11] improved this to  $C_k \leq (81/64)k^2/(\ln k)$ . Very recently Körner [4] showed  $C_k \leq (k/2) + o(1)$  (again for  $n \rightarrow \infty$ ), which is conjectured to be tight in [7]. The case  $n = 2k$  was discussed by Lovász, Nešetřil and Pultr [8], they proved that the dimension of the product of  $d$  (nontrivial) complete graphs is  $d$ . This implies  $\dim_P(K(2k, k)) = \lceil \log_2 \binom{2k}{k} \rceil = 2k - O(\log k)$ .

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The aim of this note is to point out another connection between the Prague dimension of graphs and the dimension theory of partially ordered sets by giving a very short proof of the upper bound in (1).

## 2. SCRAMBLING PERMUTATIONS AND DIMENSION OF POSETS

The dimension of a partially ordered set  $P$  is the minimum  $d$  such that  $P$  can be embedded into  $\mathbf{R}^d$  in an order preserving way. In other words, it is the minimum number of linear extensions  $\pi_1, \dots, \pi_d$  such that for all  $x, y \in P$  there exists a  $\pi_i$  with  $x <_i y$  ( $x$  precedes  $y$  in  $\pi_i$ ) except, of course, if  $y <_P x$ . In the latter case  $y$  precedes  $x$  in all linear extensions. Additional background material on dimension theory can be found in the monograph [13].

Let  $2^S$  denote the collection of subsets of  $S$ , and let  $\mathbf{B}_n = (2^{[n]}, \subseteq)$  denote the *Boolean lattice*, the subsets of  $[n]$  ordered by inclusion. For a set  $S$ , let  $\binom{S}{k}$  denote the collection of  $k$ -element subsets of  $S$ . For  $0 \leq s < t \leq n$  let  $\mathbf{B}_n(s, t)$  denote the restriction of  $\mathbf{B}_n$  to  $\binom{[n]}{s} \cup \binom{[n]}{t}$ . Finally, let  $\dim(n; s, t)$  denote the (order) dimension of  $\mathbf{B}_n(s, t)$ . The function  $\dim(n; s, t)$  was first studied by Dushnik [1] in 1950, he determined the exact value for  $\dim(n; 1, t)$  when  $2\sqrt{n} - 2 \leq t < n - 1$ .

Call the set of permutations of  $[n]$ ,  $\Pi$ ,  $t$ -*scrambling* if for every (now unordered)  $t$ -subset  $\{p_1, \dots, p_t\} \subset [n]$  and for every distinguished element of the set, say  $p_j$ , there is a permutation  $\pi \in \Pi$  such that  $\pi(p_j)$  precedes all the other  $(t-1)$   $p_i$ 's. The cardinality of the smallest  $t$ -scrambling family is denoted by  $N(n, t)$ . It is easy to see that determination of  $N(n, t)$  is equivalent to the question of the dimension of the partially ordered set formed by the  $(t-1)$  and 1-element subsets of  $[n]$  and ordered by inclusion, i.e.,  $N(n, t) = \dim(n; 1, t-1)$ . For  $t$  is fixed and  $n \rightarrow \infty$  an argument due to Hajnal and Spencer [12] gives that

$$(2) \quad \log_2 \log_2 n \leq N(n, t) \leq \frac{t}{\log_2(2^t/(2^t - 1))} \log_2 \log_2 n.$$

In [3] the asymptotic  $N(n, 3) = \log_2 \log_2 n + (\frac{1}{2} + o(1)) \log_2 \log_2 \log_2 n$  was proved.

**Theorem 2.1.**  $\dim_P(K(n, k)) \leq N(n, 2k - 1)$ .

*Proof.* Let  $\pi_1, \dots, \pi_d$  be a  $(2k - 1)$ -scrambling set of permutations of  $[n]$ . We define  $\varphi_1, \dots, \varphi_d$  good colorings of the Kneser graph  $K(n, k)$ ,  $\varphi_i : \binom{[n]}{k} \rightarrow [n]$ , as follows. Let  $\varphi_i(K) = x$  where  $x \in K$  is the smallest element of  $K$  in the linear order  $\pi_i$ .

As  $\varphi_i(K) \in K$ , for disjoint  $k$ -sets,  $K, L \in \binom{[n]}{k}$ , we have that  $\varphi_i(K) \neq \varphi_i(L)$  for all  $i$ . However, for a non-edge, i.e. for an intersecting pair  $(K, L)$ , for  $x \in K \cap L$ , one can find a permutation  $\pi_i$  which puts  $x$  to the first place among the elements in  $K \cup L$ .  $\square$

**Remark 2.2.** The constructions in [10, 11, 4] use qualitatively independent partitions and  $k$ -independent families of sets. Let us note that the upper bound in (2) also uses  $k$ -independent families of sets so it cannot give a better bound for  $C_k$  as  $2^k$ . However, together with the upper bound from [3] for  $N(n, 3)$ , it gives the asymptotic for the case  $k = 2$ , which was also showed in [10]. Finally, Theorem 2.1 also gives a number of new upper bounds

for  $\dim_P(K(n, k))$ , when  $n$  is “not too large” with respect to  $k$ , e.g.,  $k \sim \log n$ , where Kierstead’s bound [5] gives  $O(\log^3 n / \log \log n)$ .

**Remark 2.3.** One can easily see, that, similarly to the examples in [10, 11, 4], our construction is *faithful*, i.e.,  $\varphi(K) \cap \varphi(L) = K \cap L$  holds for every two  $k$ -sets, where  $\varphi(K) := \{\varphi_i(K) : 1 \leq i \leq d\}$ .

**Remark 2.4.** (*Binary intersection representations.*) Körner and Monti [6] defined the *Bohemian representation* of the Kneser graph  $K(n, k)$  as a set of colorings of its vertex set,  $\varphi_1, \dots, \varphi_t$ , where now  $\varphi_i : \binom{[n]}{k} \rightarrow N$  is not necessarily a good coloring of the graph, and a function  $\varphi : 2^{[t]} \rightarrow 2^{[n]}$  with the following property. For a pair of distinct sets  $A, B \in \binom{[n]}{k}$  let  $\delta(A, B)$  denote a sequence from  $\{0, 1\}^t$  with  $\delta_i = 1$  for  $\varphi_i(A) = \varphi_i(B)$  and 0 otherwise. In a Bohemian representation  $(\varphi_1, \dots, \varphi_t, \varphi)$  we want to be able to read out the intersection structure of the complete hypergraph knowing only the binary vectors,  $\delta(A, B)$ , i.e., we have  $\varphi(\delta(A, B)) = A \cap B$ . The minimum of such  $t$  is called the *Bohemian dimension*, and is denoted by  $T(n, k)$ . Körner and Monti [6] proved that

$$k - 1 \leq \liminf_{n \rightarrow \infty} \frac{T(n, k)}{\log_2 n} \leq \limsup_{n \rightarrow \infty} \frac{T(n, k)}{\log_2 n} \leq k(k - 1).$$

Using a different kind of set of scrambling permutations, one can see that  $T(n, k) = O(\log n)$  as  $k$  is fixed and  $n \rightarrow \infty$  as follows. Call a family of permutations  $\pi_1, \dots, \pi_t$  of  $[n]$  *completely  $k$ -scrambling* if for every ordered  $k$ -subset  $\{p_1, \dots, p_k\}$  of  $k$  distinct elements of  $[n]$  there is a permutation  $\pi_i$  with  $\pi_i(p_1) < \dots < \pi_i(p_k)$ . This means that all  $k$ -subsets appear in all  $k!$  possible orderings. The cardinality of the smallest completely  $k$ -scrambling family is denoted by  $N^*(n, k)$ . It is known (for  $k \geq 3$ ) that  $\frac{1}{2}(k - 1)! \log_2 n < N^*(n, k) \leq (1 + o(1)) \frac{k}{\log_2(k!/(k! - 1))} \log_2 n$ . Here the lower bound is from [2] and the upper bound is due to Spencer [12].

Now, one can easily see, that a completely  $(4k - 2)$ -scrambling set of permutations in the same way as in Theorem 2.1 provides a Bohemian representation of  $K(n, k)$  thus proving  $T(n, k) \leq N^*(n, 4k - 2)$ . Even more, again, the obtained  $\varphi_i$ ’s are proper colorings of the Kneser graph.

Further problems and connections between permutations and order dimensions can be found in [2].

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