

## On Automorphisms of Line-graphs

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Let  $\alpha$  be an automorphism of the line-graph of the  $r$ -uniform hypergraph  $\mathcal{H}$  with  $n$  points. If the valencies of  $\mathcal{H}$   $v(x_1) \leq v(x_2) \leq \dots \leq v(x_n)$  and

$$v(x_2) \geq v(n, r) \quad \text{and} \quad v(x_{2r}) > v(n, r) = \binom{n-1}{r-1} - \binom{n-r-1}{r-1} + 1,$$

then for  $n > 4r$   $\alpha$  is induced by an automorphism of  $\mathcal{H}$  (i.e. a permutation of  $V(\mathcal{H})$ ).

Two examples show that the valency conditions of above theorem cannot be weakened in any point of  $\mathcal{H}$ .

### 1. INTRODUCTION, RESULTS

Let  $X = \{x_1, \dots, x_n\}$  denote the points of an  $r$ -uniform hypergraph  $\mathcal{H}$ . Let  $L(\mathcal{H})$  be the line-graph of  $\mathcal{H}$ , i.e. the underlying set of  $L(\mathcal{H})$  is the edges of  $\mathcal{H}$  and the pair  $(E_1, E_2)$ ,  $E_1 \neq E_2$ ,  $E_i \in \mathcal{H}$  is an edge of  $L(\mathcal{H})$  iff  $E_1 \cap E_2 \neq \emptyset$ . Let  $\text{Aut}(L(\mathcal{H}))$  be its automorphisms. Denote the set of  $r$ -tuples of  $X$  by  $K'_n$ .

Every permutation  $\alpha \in \text{Aut } \mathcal{H}$  of  $X$  induces an automorphism  $a_\alpha$  of  $L(\mathcal{H})$  in a natural way, namely  $a_\alpha(E) = \{\alpha(x) : x \in E\}$  for every  $E \in \mathcal{H}$ .

C. Berge and J. C. Fournier proved the following theorem [1, 2].

**PROPOSITION 1.** *If  $a \in \text{Aut}(L(K'_n))$  and  $n > 2r$ , then there exists a permutation  $\alpha$  on  $X$  which induces  $a$  (i.e.  $a = a_\alpha$ ).*

The condition  $n > 2r$  cannot be omitted.

This question is strongly connected with the problem of reconstructing an  $r$ -graph  $\mathcal{H}$  from its line-graph  $L(\mathcal{H})$ . For graphs the following theorem of Whitney [5] is well known.

**PROPOSITION 2.** *If every vertex of the graph  $\mathcal{G}$  has valency  $> 3$ , then it can be reconstructed from its line-graph.*

Therefore, under the condition of Proposition 2 if  $a \in \text{Aut } L(\mathcal{G})$ , then there exists a permutation  $\alpha$  of vertices of  $\mathcal{G}$  for which  $a = a_\alpha$ .

The following theorem is an extension of Proposition 1 and generalized Proposition 2 to  $r$ -graphs. Henceforth the *valency* or *degree* of the point  $x$  of the hypergraph  $\mathcal{H}$  is, as usual,  $v(x) = |\{E : x \in E \in \mathcal{H}\}|$ .

**THEOREM 1.** *If every vertex of the  $r$ -uniform hypergraph  $\mathcal{H}$  has valency greater than*

$$v(n, r) = \binom{n-1}{r-1} - \binom{n-r-1}{r-1} + 1$$

*and  $n > 2r$ , then for every  $a \in \text{Aut}(L(\mathcal{H}))$  there exists a permutation  $\alpha$  on  $X$  such that  $a = a_\alpha$ .*

For graphs, Theorem 1 seems slightly weaker than Proposition 2 but, as a matter of fact, they are equivalent. The reason for this is that  $v(n, 2) = 3$  is independent of  $n$  while  $v(n, r)$  tends to infinity for fixed  $r \geq 3$ .

The following example shows that the claim of Theorem 1 is sharp in the sense that  $v(n, r)$  cannot be replaced by a smaller integer if  $n$  is great enough.

EXAMPLE 1. Let  $F_1 = \{x_1, \dots, x_r\}$ ,  $F_2 = \{x_{r+1}, \dots, x_{2r}\}$  and  $\mathcal{H}_1 = \{F_1, F_2\} \cup \{F \subset X: |F| = r, F \cap F_1 = \emptyset \text{ iff } F \cap F_2 = \emptyset\}$ . Then

$$v(x_i) = v(n, r) \quad \text{if } i \leq 2r$$

and

$$v(x_i) = \binom{n-1}{r-1} - 2 \binom{n-r-1}{r-1} + 2 \binom{n-2r-1}{r-1} \quad \text{if } i > 2r.$$

Thus, if  $n > 2r^2$ , then  $v(x_i) > v(n, r)$  for  $i > 2r$ . Finally the following automorphism  $a_1$  of  $L(\mathcal{H}_1)$  cannot be induced by any permutation of  $X$ .

$$a_1(F) = \begin{cases} F_2 & \text{if } F = F_1 \\ F_1 & \text{if } F = F_2 \\ F & \text{otherwise.} \end{cases}$$

EXAMPLE 2. Let  $r \geq 3$  and

$$F_3 = \{x_1\} \cup \{x_3, x_4, \dots, x_{r+1}\},$$

$$F_4 = \{x_2\} \cup \{x_3, x_4, \dots, x_{r+1}\}$$

and

$$\mathcal{H}_2 = \{F_3, F_4\} \cup \{F \subset X: |F| = r, F \cap F_3 = \emptyset \text{ iff } F \cap F_4 = \emptyset\}.$$

Then

$$v(x_1) = v(x_2) = v(n, r) - 1,$$

$$v(x_3) = v(x_4) = \dots = v(x_{r+1}) = \binom{n-1}{r-1} > v(n, r),$$

and if  $i > r+1$ ,  $n \geq 3r$  then

$$v(x_i) = \binom{n-1}{r-1} - 2 \binom{n-r-2}{r-2} > v(n, r).$$

Finally the following automorphism  $a_2$  of  $L(\mathcal{H}_2)$  cannot be induced by any permutation of  $X$ .

$$a_2(F) = \begin{cases} F_4 & \text{if } F = F_3 \\ F_3 & \text{if } F = F_4 \\ F & \text{otherwise.} \end{cases}$$

As we have seen, the valency condition of Theorem 1 cannot be weakened in every point of  $\mathcal{H}$ . The following theorem shows that it can be done in fewer than  $2r$  points.

THEOREM 2. Let  $\mathcal{H}$  be an  $r$ -uniform hypergraph on  $|X| = n \geq 4r$  points, and  $v(x_1) \leq v(x_2) \leq \dots \leq v(x_n)$ . If  $v(x_2) \geq v(n, r)$  and  $v(x_{2r}) > v(n, r)$ , then for every  $a \in \text{Aut}(L(\mathcal{H}))$  there exists a permutation  $\alpha$  on  $X$  such that  $a = a_\alpha$ .

Examples 1 and 2 show that the conditions of Theorem 2 cannot be weakened in any point of  $\mathcal{H}$ .

Let a set system be called intersecting if the pairwise intersections are non-empty. The value of  $v(n, r)$  in the theorems comes from the following theorem of Hilton and Milner [3].

**THEOREM 3.** *If  $\mathcal{C}$  is an  $r$ -uniform intersecting set system on  $X$ , and  $\bigcap \mathcal{C} = \emptyset$  then*

$$|\mathcal{C}| \leq v = v(n, r) = \binom{n-1}{r-1} - \binom{n-r-1}{r-1} + 1.$$

For  $n > 2r$  and  $r \neq 3$  the equality holds iff there are a point  $x \in X$  and  $r$ -tuple  $D \subset X (x \notin D)$  for which

$$\mathcal{C} = \mathcal{C}_{x,D} = \{E: x \in E \subset X: |E| = r, E \cap D \neq \emptyset\} \cup \{D\}. \quad (1)$$

If  $n > 2r$  and  $r = 3$ , then there is another extremum

$$\mathcal{C} = \mathcal{C}_D = \{E: E \subset X, |E| = 3, |E \cap D| \geq 2\}. \quad (2)$$

## 2. PROOF OF THEOREM 1

Let us suppose that  $\mathcal{H}$  satisfies the conditions of Theorem 1. Let  $a \in \text{Aut}(L(\mathcal{H}))$ . The proof is constructive. First we define a permutation  $\alpha$  on  $X$ , then show that  $a = a_\alpha$ .

Let  $\mathcal{C}_i$  denote the system of edges of  $\mathcal{H}$  containing point  $x_i$ . Similarly  $\mathcal{C}_p$  denotes the system of edges containing  $p$ .

**LEMMA 1.** *If  $\mathcal{C}$  is an  $r$ -uniform set system on  $X$  and  $|\mathcal{C}| \geq v - 1$ , then  $|\bigcap \mathcal{C}| \leq 1$ .*

Indeed, if  $n > r$ , then

$$v - 1 = \binom{n-1}{r-1} - \binom{n-r-1}{r-1} > \binom{n-2}{r-2}.$$

**PROOF OF THEOREM 1.** Since  $a(\mathcal{C}_i) = \{a(E): E \in \mathcal{C}_i\}$  is an intersecting set system and  $|a(\mathcal{C}_i)| = |\mathcal{C}_i| > v(n, r)$ , thus, according to Theorem 3, there is a common point of sets of  $a(\mathcal{C}_i)$ . This point is unique by Lemma 1. Let us denote it by  $a(x_i)$ . Therefore if

$$x_i \in E \quad \text{then} \quad \alpha(x_i) \in a(E) \quad (3)$$

for every  $x_i \in X$ .

On the other hand if  $i \neq j$ , then  $\alpha(x_i) \neq \alpha(x_j)$ . Suppose the contrary. Then  $a(\mathcal{C}_i) \cup a(\mathcal{C}_j)$  is an intersecting set system and hence  $\mathcal{C}_i \cup \mathcal{C}_j$  is an intersecting system, too.

By Theorem 3 there is a point  $p$  of  $X$  such that  $p \in \bigcap (\mathcal{C}_i \cup \mathcal{C}_j)$ . If, e.g.  $p \neq x_i$ , then  $\{x_i, p\} \subset \bigcap \mathcal{C}_i$ , and it contradicts Lemma 1. So we proved that  $\alpha$  is a permutation of  $X$ , and, by (3),  $a = a_\alpha$ .

**REMARK.** As a matter of fact we proved the following assertion. In the class of those  $r$ -uniform hypergraphs  $(Y, \mathcal{F})$  for which  $\min_{y \in Y} v_{\mathcal{F}}(y) > v(|Y|, r)$ , the line-graphs of the hypergraphs are isomorphic iff the hypergraphs themselves are isomorphic. Since  $v(n, 2) = 3$  for every  $n$ , it is exactly Proposition 2.

Let us remark, if  $r$  is fixed and  $n$  tends to infinity, then almost all hypergraphs are in the above class.

## 3. PROOF OF THEOREM 2 (SKETCH)

The detailed proof would contain some parts easier to prove than to understand their proofs. These proofs are left to the reader.

LEMMA 2. *If the conditions of Theorem 2 are satisfied, and  $|\mathcal{C}_x| = v(n, r)$  then either*

$$|\bigcap a(\mathcal{C}_x)| = 1 \quad (4)$$

*or there exist an  $F_x \in \mathcal{C}_x$  and an  $x' \notin a(F_x)$  so that*

$$a(\mathcal{C}_x) = \mathcal{C}_{x', a(F_x)}. \quad (5)$$

This states that the case (2) in Theorem 3 cannot be realized. The lemma can be proved indirectly.

According to the valency of the points of  $X$  let us divide the set  $X$  into three disjoint parts;  $X_{<v}$ ,  $X_v$ ,  $X_{>v} = \{x \in X: v(x) \leq v(n, r)\}$ . (Naturally  $X_{<v} = \emptyset$  or  $\{x_1\}$ .) The definition of the map  $\alpha$  on the  $X \setminus X_{<v}$  is similar to definition of  $\alpha$  in the proof of Theorem 1. If  $|\mathcal{C}_x| \geq v$  and  $\bigcap a(\mathcal{C}_x) \neq \emptyset$  then let  $\alpha(x) = \bigcap a(\mathcal{C}_x)$ . The points  $x \in (X \setminus X_{<v})$  are called regular for which claim (3) is realized. In particular, every point of  $X_{>v}$  is regular. For a non-regular point  $x$  of  $X_v$  let  $\alpha(x) = x'$  as was defined in Lemma 2. By Lemma 1, if  $x, y \in X \setminus X_{<v}$ ,  $x \neq y$ , then  $\alpha(x) \neq \alpha(y)$ . Moreover, the restriction of  $\alpha$  on  $X_{>v}$  is a permutation of  $X_{>v}$ .

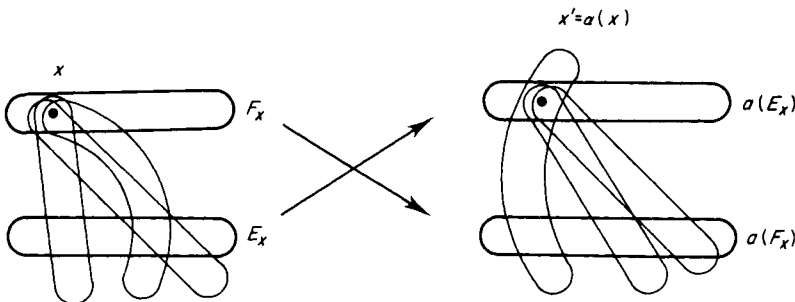
LEMMA 3. *If  $x \in X \setminus X_{<v}$  is a regular point, then*

$$x \in E \quad \text{iff} \quad \alpha(x) \in a(E). \quad (6)$$

For points of  $X_{>v}$ , (6) is evidently true. If  $x \in X_v$  is a regular point, then  $\alpha(x) \in X_v$ , so  $|a(\mathcal{C}_x)| = |\mathcal{C}_{\alpha(x)}|$ . On the other hand, by the regularity,  $a(\mathcal{C}_x) \subset \mathcal{C}_{\alpha(x)}$  so they coincide.

LEMMA 4. *If the point  $x \in X_v$  is non-regular, and  $\alpha(x) \neq x_1$ , then there exists an edge  $E_x \in \mathcal{C}_x$ , which intersects all edges in  $\mathcal{C}_x$  except for  $F_x$  defined in Lemma 2. Moreover*

- (i)  $\mathcal{C}_x = \{F_x\} \cup \{E \subseteq X: |E| = r, x \in E, E \cap E_x \neq \emptyset\}$ ,
- (ii) no point in  $E_x$  is regular (if  $n \geq 4r - 2$ ),
- (iii) the point  $\alpha(x)$  is regular.



PROOF. (i) is trivial and then  $a(\mathcal{C}_x \cup \{E_x\}) = \mathcal{C}_{x', a(F_x)} \cup \{a(E_x)\}$  for an  $x' \in a(E_x)$ . (ii) is proved indirectly. If  $y$  is regular point in  $E_x$ , let  $H \in \mathcal{C}_x$  be an edge, which contains points  $x, y$  and further  $(r-2)$  regular points, the  $\alpha$ -image of which is not in  $a(F_x)$ . Such an edge exists, because  $n \geq 4r - 2$  and there are at least  $2r - 1$  regular points of  $X$ , and there are at most  $r$  of them the  $\alpha$ -image of which is in  $a(F_x)$ . By (i),  $H \in \mathcal{C}_x$ . But  $H \neq F_x$ , therefore  $a(H) = \{\alpha(z): z \in H\}$ ,  $a(H)$  does not intersect  $a(F_x)$ . This contradicts to  $x \in H \cap F \neq \emptyset$ .

(iii) According to  $|F_x \cup E_x| = 2r$  and  $|X_{<v} \cup X_v| < 2r$ , there exists a point  $z \in (F_x \cup E_x) \cap X_{>v}$ , and, by (ii),  $z \in F_x$ . So  $\alpha(z) \in a(F_x) \cap X_{>v}$ , because  $z$  is regular. On the other hand we show, if  $x'$  is not regular then (ii) is satisfied by  $x'$  in place of  $x$  which contradicts to  $\alpha(z) \in a(F_x)$ . Indeed, Lemma 2 states that at least  $v$  among the images of the elements of the system  $\mathcal{C}_x \cup \{a(F_x)\}$  contain  $\alpha(x')$ , so  $\alpha(x') \notin X_{<v}$ .

LEMMA 5. Suppose  $v(x_1) < v$  and there exists a point  $u$  for which  $\alpha(u) = x_1$ . Let  $\alpha(x_1) = X \setminus \{\alpha(z) : z \neq x_1\}$ . Then the cycle of the permutation  $\alpha$  containing  $x_1$ , i.e.  $(x_1, \alpha(x_1), \dots, \alpha^k(x_1) = u)$ , has regular points only except possibility for  $x_1$  and  $u$ .

The proof of this lemma is similar to the proof of the Lemma 4(iii). It is not given here.

#### PROOF OF THEOREM 2

1. First we show that there are at most  $r$  non-regular points. Indeed, if there is no  $u \in X_v$  for which  $\alpha(u) = x_1$ , then  $\alpha$  is a permutation of  $X - X_{<v}$ . Let us look at the cycles of  $\alpha$  in  $X_v$ . Lemma 4(iii) claims that at least the half of the points are regular in every cycle. Therefore the number of non-regular points is  $\leq |X_{<v}| + \frac{1}{2}|X_v| \leq r$ .

If there exists  $u \in X_v$  for which  $\alpha(u) = x_1$  then, by (iii) of Lemma 4 and by Lemma 5, the number of non-regular points is at most

$$2 + \left\lceil \frac{|X_{<v} \cup X_v| - 2}{2} \right\rceil \leq 2 + \left\lceil \frac{2r - 3}{2} \right\rceil = r.$$

2. After these the application of Lemma 4(ii) gives that every  $x \in X_v$ ,  $\alpha(x) \neq x_1$  is a regular point.

3. Now it can be shown that there is no  $u \in X_v$  for which  $\alpha(u) = x_1$ .

4. It was proven that every point of  $X - X_{<v}$  is regular and  $\alpha$  is a permutation of  $X \setminus X_{<v}$ . If  $X_{<v} = \emptyset$  then  $a = a_\alpha$ , by Lemma 3. If  $v(x_1) < v$ , then let  $\alpha(x_1) = x_1$ . In this case, if  $x_1 \notin E$  then every point of  $E$  is regular, so  $a(E) = \{\alpha(x) : x \in E\}$ . Finally, if  $x_1 \in E$  then, according to Lemma 3,  $a(E)$  contains at most  $r - 1$  regular points, so  $x_1 \in a(E)$  and  $a = a_\alpha$ .

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