

Free Boolean algebras with closure operators and a conjecture of Henkin, Monk, and Tarski*

by

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Abstract

We characterize the finite-dimensional elements of a free cylindric algebra. This solves Problem 2.10 in [Henkin, Monk, Tarski: Cylindric Algebras, North-Holland, 1971 and 1985]. We generalize the characterization to quasi-varieties of Boolean algebras with operators in place of cylindric algebras.

Free algebras play an important role in universal algebra, see e.g. Andréka-Jónsson-Németi [2]. Free algebras are even more important in algebraic logic, because they give information on proof-theoretic properties of a logic. Cf. e.g. [7, §5.6], [4], [12], [8].

Cylindric algebras are special Boolean algebras with closure operators. Tarski proved that any element of a free cylindric algebra behaves the same way for all but finitely many of these operators: x is either closed for all but finitely many or only finitely many operators. See Theorem 2.6.23 in [6]. It remained open which elements of a free algebra are closed to all but finitely many, and which are closed only to finitely many. [6, 2.6.24] contains a conjecture for a characterization, but Henkin, Monk, and Tarski write there that they were unable to verify this conjecture. They also formulate the conjecture as Problem 2.10 in [6].

In this paper we solve Problem 2.10 of [6] by showing that their conjecture was right. We prove more: we give information about which element is closed under which operators, and also we generalize the result to quasi-varieties of Boolean algebras with operators.

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Boolean algebras with operators (Bo_α 's) were introduced in Jónsson-Tarski [9] and they have been well investigated ever since. See e.g. the papers in [3] and in [13], or [1], [11], [10]. Bo_α is often called also BAO in the literature. An α -dimensional *Boolean algebra with operators* is an algebra $\mathfrak{A} = \langle A, +, -, c_i, d_{ij} \rangle_{i,j \in \alpha}$ in which each $c_i : A \rightarrow A$ is a closure operator in the usual sense, and the d_{ij} 's are constants. I.e. c_i is additive and $x \leq c_i x = c_i c_i x$ holds in \mathfrak{A} . We say that \mathfrak{A} has *complemented closure operators* if the complement of a c_i -closed element is c_i -closed again, that is if $\mathfrak{A} \models \{c_i - c_i x = -c_i x : i \in \alpha\}$. An operator c_i is *normal* if $c_i 0 = 0$, and a Bo_α is normal if all of its operators are normal. α -dimensional *cylindric algebras*, CA_α 's, are Bo_α 's with commuting complemented closure operators in which three additional equations are postulated for the constants d_{ij} .

Some terminology: Let $\mathfrak{A} = \langle A, +, -, c_i, d_{ij} \rangle_{i,j \in \alpha}$ be a Bo_α and let $a \in A$. We say that a is *sensitive to an operator* c_i if $a \neq c_i a$ in \mathfrak{A} . We say that a is *finite (cofinite) dimensional* if a is sensitive to finitely many (co-finitely many) of the operators c_i ($i \in \alpha$). We say that a is *zero-dimensional* if a is not sensitive to any of the operators, i.e. if $a = c_i a$ for all $i \in \alpha$. We say that a is a *constant* element in \mathfrak{A} if a is generated by the constants $0, 1, d_{ij}$ ($i, j \in \alpha$), or in other words, if a is an element of the smallest subalgebra of \mathfrak{A} . If a generator set G of \mathfrak{A} is fixed, then by a *generating term* of a we understand a term $\tau(x_1, \dots, x_n)$ in the language of Bo_α such that $a = \tau(g_1, \dots, g_n)$ in \mathfrak{A} for some $g_1, \dots, g_n \in G$.

Theorem 1 (Solution of Problem 2.10 in [6]) *Let α be an infinite ordinal and β be a nonzero cardinal.*

- (i) *The set of all finite-dimensional elements of the β -generated free algebra $\mathfrak{F}\tau_\beta CA_\alpha$ is the subuniverse generated by $\{g \cdot -c_0(-d_{01}) : g \in G\}$, where G is the set of free generators of $\mathfrak{F}\tau_\beta CA_\alpha$.*
- (ii) *If an element is not finite-dimensional in $\mathfrak{F}\tau_\beta CA_\alpha$, then it is sensitive to all operators which do not occur in all of its generating terms.*

We will prove Theorem 1 as a corollary of the following theorem about all quasi-varieties of normal Bo_α 's.

Theorem 2 *Assume $\alpha \geq \omega$. Let K be any quasi-variety of normal Bo_α 's in which the equations (1) – (4) below are valid for all $i, j, k \in \alpha$.*

- (1) $c_i(c_i x \cdot y) = c_i c_i x \cdot c_i y$ (weak distributivity)
- (2) $c_k d_{ij} = d_{ij}$ if $k \notin \{i, j\}$

$$(3) \quad c_i d_{ij} = 1$$

$$(4) \quad c_i - d_{ij} = 1.$$

Let $\beta \neq 0$. Then (i) and (ii) below hold.

- (i) All non-constant elements of the β -generated K -free algebra $\mathfrak{F}\mathfrak{r}_\beta K$ are cofinite-dimensional.
- (ii) Moreover, a non-constant element a of the free algebra can be insensitive to an operator c_i only if c_i occurs in all generating terms of a from the free generators G of $\mathfrak{F}\mathfrak{r}_\beta K$. In more detail: Assume that $a = \tau(\bar{g})$ for some sequence \bar{g} of free generators and c_i does not occur in the term τ . Then $a \neq c_i a$ in $\mathfrak{F}\mathfrak{r}_\beta K$.

In proving Theorem 2, we will use the following two lemmas. Lemma 3 below is a version of the known theorem in cylindric algebra theory that relativization with zero-dimensional elements is a homomorphism. We have to prove this because we want to use the statement of this known theorem under much weaker conditions.

Lemma 3 *Assume that Γ is a set and \mathfrak{A} is a Bo_Γ enriched perhaps with additional constants, and assume that the equations (1) (weak distributivity) hold in \mathfrak{A} . Let b be a zero-dimensional element of \mathfrak{A} , a_1, \dots, a_n arbitrary elements of \mathfrak{A} , and let $\tau(x_1, \dots, x_n)$ be a term in the language of \mathfrak{A} . Then*

$$b \cdot \tau^{\mathfrak{A}}(a_1, \dots, a_n) = b \cdot \tau^{\mathfrak{A}}(a_1 \cdot b, \dots, a_n \cdot b).$$

Proof. We proceed by induction on τ . The only non-trivial cases are when τ is $c_i \sigma$ or $-\sigma$. In the first case we use weak distributivity, and the proof of the second case is much the same as in the cylindric algebra case. We denote $\sigma^{\mathfrak{A}}(a_1, \dots, a_n)$ and $\sigma^{\mathfrak{A}}(a_1 \cdot b, \dots, a_n \cdot b)$ by $\sigma(\bar{a})$ and $\sigma(\bar{a} \cdot b)$ respectively. Proof of homomorphism with respect to c_i : $b \cdot c_i \sigma(\bar{a}) = c_i(b \cdot \sigma(\bar{a})) = c_i(b \cdot \sigma(\bar{a} \cdot b)) = b \cdot c_i \sigma(\bar{a} \cdot b)$, using (1), $b = c_i b$ and the inductive hypothesis on σ . Proof of homomorphism with respect to complementation $-$: $b \cdot -\sigma(\bar{a}) = b \cdot -(b \cdot \sigma(\bar{a})) = b \cdot -(b \cdot \sigma(\bar{a} \cdot b)) = b \cdot -\sigma(\bar{a} \cdot b)$. ■

Lemma 4 *Assume that K satisfies the hypotheses of Theorem 2. Then the following is valid in K :*

$$c_k - d_{ij} = -d_{ij} \quad \text{if } k \notin \{i, j\}.$$

Proof. It is enough to prove that $c_k(-d_{ij}) + d_{ij} = 1$ and $c_k(-d_{ij}) \cdot d_{ij} = 0$. We will use (1) – (4), additivity and normality of the operators c_i , and of course Boolean algebra. Proof of the first equation: $c_k(-d_{ij}) + d_{ij} = c_k(-d_{ij}) + c_k(d_{ij}) = c_k(-d_{ij} + d_{ij}) = c_k 1 = c_k(d_{ki} + -d_{ki}) = c_k d_{ki} + c_k - d_{ki} = 1$. Proof of the second equation: $c_k(-d_{ij}) \cdot d_{ij} = c_k(-d_{ij}) \cdot c_k c_k d_{ij} = c_k(-d_{ij} \cdot c_k d_{ij}) = c_k(-d_{ij} \cdot d_{ij}) = c_k 0 = 0$. ■

We are ready to prove Theorem 2.

Proof of Theorem 2: It is enough to prove the second statement of Theorem 2, because it is a stronger version of the first one. Let α, β, K be as in the hypotheses of the theorem. Set $\mathfrak{F} \stackrel{\text{def}}{=} \mathfrak{F}\mathfrak{r}_\beta K$, let a be a non-constant element of \mathfrak{F} , let $\tau(x_1, \dots, x_n)$ be a term in the language of Bo_α and g_1, \dots, g_n be free generators in \mathfrak{F} such that $a = \tau(g_1, \dots, g_n)$. Let Γ be the set of all indices of operators (i.e. c_i 's) occurring in τ . We will show that $a \neq c_i a$ in \mathfrak{F} for all $i \in \alpha \setminus \Gamma$.

Let $i \in \alpha \setminus \Gamma$. If $K \not\models c_i \tau(\bar{x}) = \tau(\bar{x})$, then $c_i \tau(\bar{g}) \neq \tau(\bar{g})$ because the g_k 's are free generators in $\mathfrak{F} = \mathfrak{F}\mathfrak{r}_\beta K$, i.e. $a \neq c_i a$ and we are done.

Assume therefore that $K \models c_i \tau(\bar{x}) = \tau(\bar{x})$. We will derive a contradiction. Let $j \in \alpha \setminus \Gamma, j \neq i$. Such a j exists because α is infinite and Γ is finite. Now

$$(\star) \quad K \models \tau(\bar{x}) = c_i(\tau(\bar{x}) \cdot d_{ij}) = c_i(\tau(\bar{x}) \cdot -d_{ij})$$

by (1), (3) and (4). For every $1 \leq k \leq n$ define

$$e_k \stackrel{\text{def}}{=} g_k \cdot d_{ij} + d_{01} \cdot -d_{ij}.$$

Let \mathfrak{N} denote the Γ -reduct of \mathfrak{F} together with all the constants, i.e. the universe of \mathfrak{N} is the same as that of \mathfrak{F} and the operations of \mathfrak{N} are those of \mathfrak{F} except that we omit those operations c_k where $k \notin \Gamma$. Then \mathfrak{N} is a Bo_Γ with additional constants in which (1) holds. Also, d_{ij} is zero-dimensional in \mathfrak{N} by (2) and $i, j \notin \Gamma$. Similarly, $-d_{ij}$ is zero-dimensional in \mathfrak{N} by Lemma 4, and the term τ is in the language of \mathfrak{N} . Thus we can apply Lemma 3.

By applying (\star) and Lemma 3 twice, and using the definition of \bar{e} we get $\tau(\bar{e}) = c_i(\tau(\bar{e}) \cdot d_{ij}) = c_i(\tau(\bar{e} \cdot d_{ij}) \cdot d_{ij}) = c_i(\tau(\bar{g} \cdot d_{ij}) \cdot d_{ij}) = c_i(\tau(\bar{g}) \cdot d_{ij}) = \tau(\bar{g})$. Completely analogously we obtain $\tau(\bar{e}) = \tau(\bar{d}_{01})$. Thus $\tau(\bar{g}) = \tau(\bar{d}_{01})$. But this is a contradiction because $a = \tau(\bar{g})$ is a non-constant element by our assumption and since $\tau(\bar{d}_{01})$ is a constant element. ■

Corollary 5 *Let K be a variety of Boolean algebras with complemented closure operators. Then (2)-(4) of Theorem 2 imply the conclusion of Theorem 2.*

Proof. A Bo_α with complemented closure operators is always normal and weak distributivity (1) holds in it. This is easy to check. ■

We are ready to prove Theorem 1 now. We note that neither (i) nor (ii) of Theorem 2 is true for $K = CA_\alpha$. Indeed, if g is a free generator, then $a = g \cdot -c_0(-d_{01})$ is a zero-dimensional element in $\mathfrak{F}r_\beta CA_\alpha$, yet a is not a constant in $\mathfrak{F}r_\beta CA_\alpha$. Also, $a = c_2a$ because it is zero-dimensional, yet 2 does not occur in the generating term $\tau = x \cdot -c_0(-d_{01})$ of a .

Proof of Theorem 1. First we prove (ii). Let M denote the subuniverse of $\mathfrak{F}r_\beta CA_\alpha$ generated by $\{g \cdot -c_0(-d_{01}) : g \in G\}$. Assume that $a \in Fr_\beta CA_\alpha \setminus M$, $a = \tau(g_1, \dots, g_n)$, $g_1, \dots, g_n \in G$ and c_j does not occur in τ . We want to show that $a \neq c_j a$ in $\mathfrak{F}r_\beta CA_\alpha$.

Let K denote the variety of those CA_α 's in which the equation $c_0 - d_{01} = 1$ holds. Then K satisfies the hypotheses of Theorem 2. (We note that CA_α does not satisfy those hypotheses, because (4) fails for CA_α .) Assume that the set of free generators of $\mathfrak{F}r_\beta K$ is also G , and let $h : \mathfrak{F}r_\beta CA_\alpha \rightarrow \mathfrak{F}r_\beta K$ be a homomorphism which is the identity on G (i.e. $h(g) = g$ for all $g \in G$). Such a homomorphism exists because $\mathfrak{F}r_\beta K \in CA_\alpha$. We are going to show that $h(a)$ is non-constant in $\mathfrak{F}r_\beta K$.

Let $\delta \stackrel{\text{def}}{=} c_0(-d_{01})$. Then it is a cylindric algebraic theorem that δ is a zero-dimensional element. Cf. [6, 1.6.9(i)]. Let $e(x) \stackrel{\text{def}}{=} x \cdot \delta$ for any $x \in Fr_\beta CA_\alpha$. Then by Lemma 3, e is a homomorphism on $\mathfrak{F}r_\beta CA_\alpha$ in the sense of [6, Def.0.2.1, p.67], i.e. there is a unique algebra \mathfrak{A} such that e is a surjective homomorphism from $\mathfrak{F}r_\beta CA_\alpha$ onto \mathfrak{A} . Cf. [6, Thm.0.2.4]. Then $\mathfrak{A} \in CA_\alpha$ because CA_α is a variety. Then $\mathfrak{A} \in K$ by $\mathfrak{A} \models \delta = 1$. We have $A \subseteq Fr_\beta CA_\alpha$ because $A = \{a \cdot \delta : a \in Fr_\beta CA_\alpha\}$. Let h' denote the restriction of h to A . We are going to show that h' is an isomorphism between \mathfrak{A} and $\mathfrak{F}r_\beta K$. It is easy to check that $h' : \mathfrak{A} \rightarrow \mathfrak{F}r_\beta K$ is a homomorphism because $h(\delta) = 1$. It remains to show that h' is one-to-one on A . Let $f : \mathfrak{F}r_\beta K \rightarrow \mathfrak{A}$ be a homomorphism such that $f(g) = g \cdot \delta$ for all $g \in G$. Such a homomorphism exists, because $\mathfrak{A} \in K$ and $g \cdot \delta \in A$ for all $g \in G$. Thus the two homomorphisms $f \circ h'$ and $h' \circ f$ are homomorphisms on \mathfrak{A} and $\mathfrak{F}r_\beta K$ respectively, such that they are identity on the generating sets $\{g \cdot \delta : g \in G\}$ and G respectively. Thus both $f \circ h'$ and $h' \circ f$ are identity on the corresponding algebras, showing that h' and f both are isomorphisms.

Now we are ready to show that $h(a)$ is non-constant in $\mathfrak{F}r_\beta K$. Assume the contrary, i.e. that $h(a) = \sigma$ in $\mathfrak{F}r_\beta K$ for some constant term σ . Then $h(\sigma \cdot \delta) = h(\sigma^{\mathfrak{A}}) = h'(\sigma^{\mathfrak{A}}) = \sigma$ because h' is a homomorphism, and then $a \cdot \delta = \sigma \cdot \delta$ because h is one-to-one on A (since h' is an isomorphism) and $a \cdot \delta \in A$. Since δ is zero-dimensional, so is $-\delta$, and thus $a \cdot -\delta = \tau(\bar{g} \cdot -\delta)$ by Lemma 3. So $(a \cdot -\delta) \in M$ by

the definition of M . This is a contradiction, because $a \notin M$ but we have seen that $a \cdot \delta = \sigma \cdot \delta$ is a constant, hence in M , and M is closed under addition. We have shown that $h(a)$ is not a constant in $\mathfrak{Fr}_\beta K$.

Now we can apply Theorem 2(ii), which yields that $h(a) \neq c_j h(a)$ in $\mathfrak{Fr}_\beta K$ because τ is also a generating term for $h(a)$ in $\mathfrak{Fr}_\beta K$. But then $a \neq c_j a$ in $\mathfrak{Fr}_\beta CA_\alpha$ because h is a homomorphism. This completes the proof of (ii).

To prove (i), it is enough now to show that each element of M is finite-dimensional. This follows from the cylindric algebraic theorem that $\eta \stackrel{\text{def}}{=} -c_0(-d_{01})$ is a hereditarily zero-dimensional element, which means that not only η , but all elements smaller than η are zero-dimensional. See [6, 1.6.20]. Thus $g \cdot \eta$ is zero-dimensional for each $g \in G$, and then one can apply the easy cylindric algebraic theorem that finite-dimensional elements generate finite-dimensional ones (cf. [6, 2.1.5(i)]). ■

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