

AN EXTREMAL PROBLEM FOR COMPLETE BIPARTITE GRAPHS

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Dedicated to the memory of Paul Turán

Abstract

Define $f(n, k)$ to be the largest integer q such that for every graph G of order n and size q , \bar{G} contains every complete bipartite graph $K_{a,b}$ with $a+b=n-k$. We obtain (i) exact values for $f(n, 0)$ and $f(n, 1)$, (ii) upper and lower bounds for $f(n, k)$ when $k \geq 2$ is fixed and n is large, and (iii) an upper bound for $f(n, \lfloor \varepsilon n \rfloor)$.

1. Introduction

Extremal graph theory, which was initiated by Turán in 1941 [4], is still the source of many interesting and difficult problems. The standard problem is to determine $f(n, G)$, the smallest integer q such that every graph with n vertices and q edges contains a subgraph isomorphic to G . It is striking that whereas Turán completely determined $f(n, K_m)$, there is much which is as yet unknown concerning $f(n, K_{a,b})$. In this paper, we consider a variant of the extremal problem for complete bipartite graphs. In this variant we ask how many edges must be deleted from K_n so that the resulting graph no longer contains $K_{a,b}$ for some pair (a, b) with $a+b=m$. Specifically, we seek to determine an extremal function $f(n, k)$ defined as follows. For $m > 1$, let B_m denote the class of all graphs G such that $G \supset K_{a,b}$ for every pair (a, b) with $a+b=m$. Then for $n > k+1$, $f(n, k)$ is the largest integer q such that every graph G of order n and size $\binom{n}{2} - q$ is a member of B_{n-k} . In this paper we obtain exact values for $f(n, 0)$ and $f(n, 1)$, upper and lower bounds for $f(n, k)$ when $k > 1$ is fixed and n is large, and an upper bound for $f(n, \lfloor \varepsilon n \rfloor)$.

2. Terminology and notation

All graphs considered in this paper will be ordinary graphs, i.e. finite, undirected graphs, without loops or multiple edges.

A graph with vertex set V and edge set E will be denoted $G(V, E)$. If $|V|=p$ and $|E|=q$, G is said to be of order p and size q . With $X, Y \subseteq V$, the set of edges in E of the form $\{x, y\}$ where $x \in X$ and $y \in Y$ will be denoted $E(X, Y)$. The complement of G will be denoted \bar{G} .

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The size of G will be given by $q(G)$. The order of the *largest connected component* will be given by $\mu(G)$ and the order of the *smallest connected component* will be given by $\eta(G)$. In particular, $\eta(G)=1$ means that G contains an isolated vertex.

Let A be a finite set. Then A^k will denote the Cartesian product $A \times A \times \dots \times A$ with k factors and $[A]^k$ will denote the collection of k -element subsets of A .

Where x is a real number, $[x]$ and $\lceil x \rceil$ denote the greatest integer $\leq x$ and the least integer $\geq x$, respectively.

For any notation or terminology not explicitly mentioned in this section, we refer the reader to [1] or [2].

3. Calculation of $f(n, k)$ where k is fixed

Our starting point is the following simple observation. If G is of order n and $\mu(G) > \lfloor n/2 \rfloor$, then $\bar{G} \not\in K_{a,b}$ with $a = \lfloor n/2 \rfloor$, $b = \lfloor n/2 \rfloor$ and so $\bar{G} \notin B_n$. The opposite direction is described by the following useful lemma.

LEMMA 1. *If $G(V, E)$ is a graph of order n such that (i) $\mu(G) \leq \lfloor n/2 \rfloor$, (ii) $\eta(G) = 1$, and (iii) $q(G) \leq \lfloor 2n/3 \rfloor - 1$, then $\bar{G} \in B_n$. This result is sharp.*

PROOF. The proof is by induction on n . If $n=2$, then G is required to be empty and so the conclusion holds. Let $\mu(G)=k$. It is easy to see that the result holds if $k=1$ or 2, so we may assume that $k \geq 3$. Let $H=G-X$, where X is a component of order k . Then H is a graph of order $n-k$ and $\eta(H)=1$. Now $q(H) \leq \lfloor 2n/3 \rfloor - k \leq \lfloor 2(n-k)/3 \rfloor - 1$, the second inequality being by virtue of the fact that $k \geq 3$. Also, $\mu(H) \leq \min(k, \lfloor 2n/3 \rfloor - k + 1)$. If $3k \leq n$, then $k \leq \lfloor (n-k)/2 \rfloor$ and if $3k \geq n+1$, then $\lfloor 2n/3 \rfloor - k + 1 \leq \lfloor (n-k)/2 \rfloor$. Hence, in all cases H satisfies (i)–(iii) and so, by the induction hypothesis, $\bar{H} \in B_{n-k}$. Since X and \bar{H} are completely joined in \bar{G} , it follows that $\bar{G} \in B_n$.

From the remark made earlier, we know that condition (i) cannot be weakened. To see that (ii) cannot be weakened, note that if $\eta(G) > 1$, then $\bar{G} \not\in K_{1, n-1}$. Finally, with $n \geq 7$ set $m = \lfloor (n+1)/3 \rfloor + 1$, $k = \lfloor n/3 \rfloor + 1$, $l = n - m - k$ and consider the graph $G = T_m \cup T_k \cup \bar{K}_l$, where T_m and T_k denote arbitrary trees of orders m and k , respectively. In this case, we have $\mu(G) \leq \lfloor n/2 \rfloor$, $\eta(G) = 1$ and $q(G) = \lfloor 2n/3 \rfloor$. However, $\bar{G} \not\in K_{a,b}$ with $a = \lfloor 2n/3 \rfloor + 1$, $b = \lfloor n/3 \rfloor - 1$. This example shows that condition (iii) cannot be weakened. \square

With the aid of Lemma 1, we can obtain the exact value of $f(n, k)$ in case $k=0$ or 1.

THEOREM 1. *For all $n \geq 2$, $f(n, 0) = \lfloor n/2 \rfloor - 1$ and for all $n \geq 3$, $f(n, 1) = \lfloor (n+1)/2 \rfloor$.*

PROOF. With $m = \lfloor n/2 \rfloor + 1$, let $G = T_m \cup \bar{K}_{n-m}$, where T_m denotes an arbitrary tree of order m . Thus, G is a graph of order n , $q(G) = \lfloor n/2 \rfloor$ and $\mu(G) = \lfloor n/2 \rfloor + 1$. Since $\mu(G) > \lfloor n/2 \rfloor$, it follows that $\bar{G} \notin B_n$ and this example shows that $f(n, 0) \leq \lfloor n/2 \rfloor - 1$. To prove the inequality in the other sense, consider an arbitrary graph G of order n and size $q(G) \leq \lfloor n/2 \rfloor - 1$. Note that such a graph must satisfy (i) $\mu(G) \leq \lfloor n/2 \rfloor$, (ii) $\eta(G) = 1$, and (iii) $q(G) \leq \lfloor 2n/3 \rfloor - 1$. Hence, by Lemma 1, $\bar{G} \in B_n$.

With $m = \lfloor (n+1)/2 \rfloor + 1$, let $G = C_m \cup \bar{K}_{n-m}$, where C_m denotes the cycle of order m . Thus, G is a graph of order n and size $q(G) = \lfloor (n+1)/2 \rfloor + 1$. Moreover,

if x is an arbitrary vertex of G , then $\mu(G-x) \cong [(n+1)/2] > [(n-1)/2]$. It follows that for each $x \in \overline{G-x} \cong K_{a,b}$ with $a = [(n-1)/2]$, $b = \lfloor (n-1)/2 \rfloor$ and so this example shows that $f(n, 1) \cong [(n+1)/2]$. To prove the inequality in the other sense, consider an arbitrary graph G of order n and size $q(G) \cong [(n+1)/2]$. Let x be a vertex of maximal degree in G , and let $H = G - x$. If x has degree $\cong 2$, then $q(H) \cong \cong [(n+1)/2] - 2 = \lfloor (n-1)/2 \rfloor - 1$. If x has degree $\cong 1$, then G is the union of a collection of disjoint edges and so in this case as well $q(H) \cong \lfloor (n-1)/2 \rfloor - 1$. Therefore, by the first part of this theorem, $\overline{H} \in B_{n-1}$ and so $\overline{G} \in B_{n-1}$. \square

COROLLARY. Let $t(n)$ denote the largest integer q such that for every graph G of order n and size q , \overline{G} contains every tree of order n . For all $n \cong 2$, $t(n) = \lfloor n/2 \rfloor - 1$.

PROOF. Since each tree of order n is contained in an appropriate complete bipartite graph $K_{a,b}$ with $a+b=n$, it follows that $t(n) \cong f(n, 0) = \lfloor n/2 \rfloor - 1$. On the other hand, the graph $G = (n/2)P_2$ (n even) or $G = ((n-3)/2)P_2 \cup P_3$ (n odd) is a graph of order n and size $q(G) = \lfloor n/2 \rfloor$ such that $\overline{G} \cong K_{1, n-1}$. (Here, mH is used to denote the graph with m components, each isomorphic to H .) This example shows that $t(n) \cong \lfloor n/2 \rfloor - 1$. \square

At this point, one may be tempted to conjecture that for each fixed value of k , $f(n, k) = n/2 + O(1)$, perhaps even exactly calculable as in the case of $k=0$ or $k=1$. In fact, we find that for all $k \cong 2$, $n/2 + A\sqrt{n} < f(n, k) < n/2 + B\sqrt{n}$, where the positive numbers A and B depend only on k . Thus, there is a very striking difference between the case of $k=1$ and that of $k=2$. In order to establish the facts concerning the behavior of $f(n, k)$ when $k \cong 2$, we shall need several preliminary results.

The following lemma uses the term *suspended path*. A path x_0, x_1, \dots, x_k in a graph G will be called suspended if its interior vertices x_1, \dots, x_{k-1} are of degree 2 in G , whereas its end vertices (x_0 and x_k) have degree $\neq 2$.

LEMMA 2. Any tree having k vertices of degree 1 is the union of at most $2k-3$ edge-disjoint suspended paths.

PROOF. The proof is left to the reader.

LEMMA 3. Let T be a tree of order $n+1$ where $n \cong 2$. There exists a vertex x such that $\mu(T-x) \cong \lfloor n/2 \rfloor$. Consequently, there is a partition of the components of $T-x$ into two parts such that each part has at least $\lfloor n/3 \rfloor$ vertices.

PROOF. The proof is left to the reader.

LEMMA 4. Let $G(V, E)$ be a connected graph of order p and size $p+l-1$. With $k \cong 2$, set $\delta = \min(\lfloor k/2 \rfloor / (4l-3), 1/4)$. Then, there exists $X \in [V]^k$ such that $\mu(G-X) \cong \lfloor (1-\delta)p \rfloor$.

PROOF. Delete l edges from G in such a way that the resulting graph H is still connected, i.e. so that H is a tree. The deleted edges determine a subtree T in the following way. First, we find those vertices which were incident in G with one of the deleted edges and so define a set A . Then, we define T to be the union of all paths in H which join pairs of vertices from A . Let A_1 denote the vertices of A which have degree 1 in T and set $A_2 = A - A_1$. According to Lemma 2, T is the union of

at most $2|A_1|-3$ edge-disjoint suspended paths. The vertices of A_2 now subdivide these suspended paths into what we shall call *elementary paths*. The elementary paths may be described in the following way. The end-vertices of the elementary paths are precisely those vertices x such that either (i) $x \in A$ or (ii) $\deg(x) > 2$ in T . Suppose that there are r elementary paths P_1, P_2, \dots, P_r . Since $|A| \leq 2l$, it follows that $r \leq 2|A_1| + |A_2| - 3 \leq 4l - 3$.

Note the following useful property of the construction described thus far. Suppose that x is a vertex of G and that it is not a vertex of T . Then, there is a unique path in G from x to T . If there were two such paths, then one of them would have to use one of the edges which were deleted in going from G to H . This would put x on a path in H joining two vertices from A and so force x to belong to T . In light of this property, we note that the collection of elementary paths P_1, P_2, \dots, P_r may be used to define a partition $V = (V_1, V_2, \dots, V_r)$ of the vertices of G according to the following scheme. If x is an end-vertex of one or more elementary paths, it is identified with an arbitrarily chosen one of those paths. If x is an interior vertex of an elementary path, it is identified with that path. Finally, if x is a vertex of G which is not a vertex of T , let w be the other end-vertex of the unique path from x to T and identify x with the same elementary path as is w .

Now we are ready to describe and put to use the crucial properties of the elementary paths. Let u_i and v_i be the end-vertices of the i th elementary path, P_i . Our construction insures that if x is any vertex of V_i other than u_i or v_i , every path from x to a vertex in $V - V_i$ contains either u_i or v_i . In other words, by deleting u_i and v_i from G , we completely disconnect the vertices of V_i from the remaining vertices of G . Without loss of generality, we may suppose that $|V_1| \geq \dots \geq |V_r|$. Set $m = \min\{\lfloor r/4 \rfloor, \lfloor k/2 \rfloor\}$ and consider the graph $G - X$, where $X = \{u_i, v_i, i = 1, \dots, m\}$. Since $|V_1| + \dots + |V_m| \geq mp/r \geq \delta p$, it follows that $\mu(G - X)$ satisfies the stated bound unless $|V_1| > [(1 - \delta)p]$. In case $|V_1| > [(1 - \delta)p]$, set $B = V_1 \cup \{u_1, v_1\}$ and consider the tree T' spanned by the vertices of B . By Lemma 3, there exists a vertex x of this tree such that the components of $T' - x$ can be partitioned into two parts, each of cardinality at least $\lfloor (|V_1| - 1)/3 \rfloor$. Now we may delete x and either u_1 or v_1 , whichever is appropriate, and so disconnect from G a set of at least $\lfloor p/4 \rfloor$ vertices. In this case, for $X = \{x, u_1\}$ or $\{x, v_1\}$ we obtain $\mu(G - X) \leq \lfloor 3p/4 \rfloor$. \square

Now we are prepared to prove our theorem concerning $f(n, k)$ with $k \geq 2$.

THEOREM 2. *Let $k > 1$ be fixed and set $A = \sqrt{\lfloor k/2 \rfloor / 16}$ and $B = \sqrt{3k(k-1)/(k+1)}$. Then, for all sufficiently large n ,*

$$n/2 + A\sqrt{n} < f(n, k) < n/2 + B\sqrt{n}.$$

PROOF. Let $G(V, E)$ be a graph of order n and size $q = n/2 + \Delta$, where $\Delta = A\sqrt{n}$. We wish to prove that there exists $X \subseteq [V]^k$ such that $G - X$ satisfies the conditions of Lemma 1. This will establish the lower bound for $f(n, k)$. Since $\Delta = o(n)$, it follows that the number of connected components of G is at least $n - q = n/2 - o(n)$. Consequently, $\eta(G) \leq 2$. On the other hand, if $\eta(G) = 2$, then $\mu(G) = o(n)$ and so by deleting just one vertex from G we obtain a graph which satisfies the conditions of Lemma 1. Hence, we now assume that $\eta(G) = 1$. Since this is the case, we may assume that $\mu(G) > \lfloor (n-k)/2 \rfloor$, in fact $\mu(G) > \lfloor (n+k)/2 \rfloor$ for, otherwise, we may

simply delete any k vertices from the largest component. Suppose that the largest component is of order p and size $p+l-1$. Hence, we have the bounds $p \leq q = n/2 + \Delta$ and $l \leq q - p + 1 \leq \Delta$. With a view toward applying Lemma 4, note that if $\delta = \lfloor k/2 \rfloor / (4l - 3)$ then $(1 - \delta)p < (1 - \lfloor k/2 \rfloor / 4\Delta)(n/2 + \Delta) < n/2 + (\Delta^2 - \lfloor k/2 \rfloor n/8) / \Delta$. Therefore, in this case and with our choice of Δ , we have $\lfloor (1 - \delta)p \rfloor \leq \lfloor (n - k)/2 \rfloor$. Certainly if $\delta = 1/4$, $\lfloor (1 - \delta)p \rfloor \leq \lfloor (n - k)/2 \rfloor$ and so the desired result follows from Lemma 4.

The upper bound is established by the following simple construction. With m chosen to be an even integer, let H be a graph of order m which is regular of degree $k+1$ and $(k+1)$ -connected. An example of such a graph has vertices $0, 1, \dots, m-1$ with two vertices i and j joined if $i - \lfloor (k+1)/2 \rfloor \leq j \leq i + \lfloor (k+1)/2 \rfloor \pmod{m}$ and, if $k+1$ is odd, i is joined to $i + m/2$ for $1 \leq i \leq m/2$. The fact that such a graph is, indeed, $(k+1)$ -connected was proved by Harary in [3] and the proof is also given in [1, pp. 48–49]. Set $r = m(k+1)/2$ and let the edges of H be e_1, e_2, \dots, e_r . For $i = 1, 2, \dots, r$, insert a vertex y_i subdividing e_i and make y_i adjacent to $l_i - 1$ new vertices. Finally, add isolated vertices so that the resulting graph $G(V, E)$ is of order n . Thus, G is of size $q(G) = r + (l_1 + \dots + l_r)$. Without loss of generality, we may assume that $l_1 \geq l_2 \geq \dots \geq l_r$. Now make the following choices for the parameters of G . Set $m = 2 \lceil \sqrt{5kn/8(k^2 - 1)} \rceil$ and $l_1 = \dots = l_k = \lceil \sqrt{5(k-1)n/8k(k+1)} \rceil \doteq l$. Then choose l_{k+1}, \dots, l_r so that $m + (l_{k+1} + \dots + l_r) = \lfloor (n - k)/2 \rfloor + 1$. Let $Y = \{y_1, \dots, y_k\}$. It is apparent that for every $X \in [V]^k$, we have $\mu(G - X) \geq \mu(G - Y) = \lfloor (n - k)/2 \rfloor + 1$. Also, we have $q(G) = \lfloor (n - k)/2 \rfloor + 1 + kl + (k - 1)m/2 < n/2 + B\sqrt{n}$ for every $\varepsilon > 0$. Since $\mu(G - X) > \lfloor (n - k)/2 \rfloor$ for every $X \in [V]^k$, it follows that $\bar{G} \not\subseteq K_{n-k}$. This establishes the upper bound. \square

4. An upper bound for $f(n, \lfloor \varepsilon n \rfloor)$

At present, very little is known about $f(n, k)$ when $k \rightarrow \infty$ with n . However, the results of the preceding section suggest that $f(n, \lfloor \varepsilon n \rfloor) < \lfloor (1/2 + \delta)n \rfloor$ where $\delta \downarrow 0$ with ε and this much can be proved without difficulty.

THEOREM 3. Let $0 < \varepsilon < e^{-4}$ be fixed and set $\delta = \sqrt{6\varepsilon \log(1/\varepsilon)}$. For all sufficiently large values of n ,

$$f(n, \lfloor \varepsilon n \rfloor) < \lfloor (1/2 + \delta)n \rfloor.$$

PROOF. Set $p = \lfloor 1/2(1 + \delta)n \rfloor$, $q = \lfloor (1/2 + \delta)n \rfloor$, $k = \lfloor \varepsilon n \rfloor$, $r = q - p$, $a = \lfloor (n - k)/2 \rfloor$, $b = \lfloor (n - k)/2 \rfloor$, and $c = a + p - n$. Using the probabilistic method, we shall prove the existence of a graph G of order n and size $\leq q$ such that $\bar{G} \not\subseteq K_{a,b}$. Let $V = \{1, 2, \dots, n\}$, $X = \{1, 2, \dots, p\}$ and $Y = [V]^2$. The probability space used to prove the existence of G may be described as follows. Let $\Omega = \Omega_1 \times \Omega_2$ where $\Omega_1 = X^p$ and $\Omega_2 = Y^r$. Each point in Ω is given probability $1/|\Omega|$. A typical point in Ω is $\omega = (\omega_1, \omega_2)$ where $\omega_1 = (x_1, \dots, x_p)$ and $\omega_2 = (y_1, \dots, y_r)$. Corresponding to ω there is a graph defined as follows: $\{i, j\}$ is an edge in the graph for each occurrence of $x_i = j$, $x_j = i$ or $y_k = \{i, j\}$, $k = 1, \dots, r$. It is understood that any loops and/or extra edges which may be generated by the random method are simply not included in the graph so formed. If $\bar{G} \subseteq K_{a,b}$ then for some m , $c \leq m \leq a$, there are disjoint sub-

sets of X , namely A and B with $|A|=m$ and $|B|=p-k-m$, such that $E(A, B)=\varphi$. Now for fixed A and B , consider the event $E(A, B)=\varphi$. The number of points of Ω_1 in this event is $(m+k)^m(p-m)^{p-k-m}p^k$ and the number of points of Ω_2 in this event is $\left(\binom{n}{2}-m(p-k-m)\right)^r$. Hence, we obtain the bound

$$\text{Prob}(\bar{G} \supseteq K_{a,b}) \leq \sum_{m=c}^a \binom{p}{m} \binom{p-m}{k} \frac{(m+k)^m(p-m)^{p-k-m}p^k}{p^p} \left(1 - \frac{m(p-k-m)}{\binom{n}{2}}\right)^r.$$

Using Stirling's formula and some elementary bounds, we find that each term in the sum is bounded by

$$(1+2k/n)^n (p/k)^k \left(1 - a(p-k-a)\right) / \binom{n}{2}^r.$$

Substituting the values of a, k, p and r , we find that $\text{Prob}(\bar{G} \supseteq K_{a,b}) \rightarrow 0$ as $n \rightarrow \infty$ provided that $(1+2\epsilon)((1+\delta)/2\epsilon)^p (1-(1-\epsilon)(\delta-\epsilon)/2)^{\delta/2} < 1$. A simple calculation shows this to be the case when $0 < \epsilon < e^{-4}$ and $\delta = \sqrt{6\epsilon \log(1/\epsilon)}$. \square

5. Additional problems and results

The bound for $f(n, \lfloor \epsilon n \rfloor)$ provides a satisfying tie with the results for $f(n, k)$ where k is fixed; still, it leaves us with more questions than answers. Among other things, the result shows that if $F(\epsilon) \stackrel{\text{def}}{=} \lim_{n \rightarrow \infty} f(n, \lfloor \epsilon n \rfloor)/n$ exists, then $\lim_{\epsilon \downarrow 0} F(\epsilon) = 1/2$. But, does $\lim_{n \rightarrow \infty} f(n, \lfloor \epsilon n \rfloor)/n$ exist?

PROBLEM 1. For $0 < x < 1$, does $\lim_{n \rightarrow \infty} f(n, \lfloor xn \rfloor)/n$ exist?

By a variety of simple arguments, it is possible to prove bounds of the form $F_1(x) < f(n, \lfloor xn \rfloor)/n < F_2(x)$ which hold when $0 < x < 1$ is fixed and n is sufficiently large. Hence, it is at least plausible that $\lim_{n \rightarrow \infty} f(n, \lfloor xn \rfloor)/n$ exists. As an example of an upper bound for $f(n, \lfloor xn \rfloor)/n$, we give the following argument. Starting with the complete graph K_n , we wish to remove $q = \lfloor yn \rfloor$ edges e_1, e_2, \dots, e_q in such a way that all $K_{m,m}$ subgraphs with $m = \lfloor (1-x)n/2 \rfloor$ are destroyed. Having found such a number y , we are assured that $f(n, \lfloor xn \rfloor)/n < y$. Let X_i denote the set of $K_{m,m}$ subgraphs which remain after e_i has been removed. Clearly, $|X_0| = \binom{n}{m} \binom{n-m}{m}$. At the stage of removing the edge e_{i+1} there are $|X_i|$ remaining $K_{m,m}$ subgraphs and $\binom{n}{2} - i$ remaining edges. Counting multiplicity, the remaining $K_{m,m}$ subgraphs contain $|X_i| m^2$ edges. It follows that there is an edge whose removal destroys at least $|X_i| m^2 / \binom{n}{2}$ of the subgraphs in X_i . By choosing such an edge for e_{i+1} , we obtain $|X_{i+1}| \leq |X_i| \left(1 - m^2 / \binom{n}{2}\right)$. Following such a procedure for $i=1, 2, \dots, q$, we obtain

$|X_q| < \binom{n}{m} \binom{n-m}{m} \left(1 - m^2 / \binom{n}{2}\right)^q$. An easy calculation using Stirling's formula allows us to conclude that if y is chosen so that $(1 - (1-x)^2/2)^y < ((1-x)/2)^{1-x} x^x$ and n is sufficiently large, then $|X_q| = 0$. As $x \rightarrow 0$, the upper bound for $f(n, [xn])/n$ that is obtained by this argument is quite inferior to the bound given in Theorem 3. The advantage of this argument is that it is applicable for all x satisfying $0 < x < 1$.

The second problem is not concerned with the calculation of $f(n, k)$, but is certainly related to the investigation described in this paper.

PROBLEM 2. For all $n \geq 2$, determine the largest integer $m = f(n)$ such that for every tree T of order n , $T \in B_m$.

We have obtained upper and lower bounds for $f(n)$ and these results may be published elsewhere.

Finally, we note the following generalization of the basic problem considered in this paper.

PROBLEM 3. For $r \geq 2$ and $n \geq k+r$, let $f_r(n, k)$ denote the largest integer q such that for every graph G of order n and size q , $\bar{G} \supseteq K(a_1, \dots, a_r)$ for every partition (a_1, \dots, a_r) of $n-k$ into r parts. Determine $f_r(n, k)$.

The proofs given in this paper extend naturally and easily to the study of $f_r(n, k)$. For $r \geq 3$, the induction argument used in the proof of Lemma 1 yields the following result.

LEMMA. Let $r \geq 3$. If G is a graph of order n such that (i) $\mu(G) \leq [n/r]$ and (ii) $q(G) \leq [2n/(r+1)] - 1$, then $\bar{G} \supseteq K(a_1, \dots, a_r)$ for every partition (a_1, \dots, a_r) of n .

Now we can state the following generalizations of Theorems 1, 2 and 3. The reader will find that the proofs given earlier in the paper have been so structured that they readily yield the results now stated.

THEOREM. For all $r \geq 2$ and $n \geq r$, $f_r(n, 0) = [n/r] - 1$. Except for certain exceptional cases, $f_r(n, 1) = [(n-1)/r] + 1$ holds for all $r \geq 2$ and $n \geq r+1$. The exceptional cases are $f_3(4, 1) = 1$, $f_3(6, 1) = 2$, $f_3(8, 1) = 3$ and, for $r \geq 4$, $f_r(r+1, 1) = 1$ and $f_r(r+2, 1) = f_r(r+3, 1) = 2$.

THEOREM. Let $r, k > 1$ be fixed and set $A = \sqrt{[k/2]/8r}$ and $B = \sqrt{6k(k-1)/((k+1)r)}$. Then, for all sufficiently large n ,

$$n/r + A\sqrt{n} < f_r(n, k) < n/r + B\sqrt{n}.$$

THEOREM. Let $0 < \varepsilon < e^{-4}$ be fixed and set $\delta = \sqrt{r(r+1)\varepsilon \log(1/\varepsilon)}$. For all sufficiently large values of n ,

$$f_r(n, [n\varepsilon]) < [(1/r + \delta)n].$$

Exactly as in the special case of $r=2$, the methods used in this paper provide an effective means of studying $f_r(n, k)$ only when $k \ll n$. Thus, for example, the generalization of Problem 1 to consider $f_r(n, [xn])$, $0 < x < 1$, is an important

problem about which little is known at present. With n and r fixed, $f_r(n, k)$ is defined for $0 \leq k \leq n-r$, and it is worth pointing out that in addition to the $k=0$ and $k=1$ cases, $f_r(n, k)$ is known exactly for $k=n-r$. We know that $f_r(n, n-r) = (n-r+t+1)s/2-1$, where $n=(r-1)s+t$, $0 \leq t < r-1$. This is Turán's theorem.

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