

ON SETS OF DISTANCES OF n POINTS IN EUCLIDEAN SPACE

by
P. ERDŐS

Let $[P_n^{(k)}]$ be the class of all subsets $P_n^{(k)}$ of the k dimensional space consisting of n distinct points and having diameter 1. Denote by $g_k(n, r)$ the maximum number of times a given distance r can occur among n points of a set $P_n^{(k)}$. Put

$$G_k(n) = \max_r g_k(n, r), \quad g_k(n) = g_k(n, 1)$$

(i. e.) $g_k(n)$ denotes the maximum number of times the diameter can occur as a distance among n points of k dimensional space and $G_k(n)$ denotes the maximum number of times the same distance can occur between n suitably chosen points in k dimensional space). It is well known [1] that $g_2(n) = n$ and I [2] proved that

$$(1) \quad n^{1+c/\log \log n} < G_2(n) < n^{3/2}.$$

Further I conjectured that $G_2(n) < n^{1+\epsilon}$ for every $\epsilon > 0$ if $n > n_0(\epsilon)$. VÁZSONYI conjectured that $g_3(n) = 2n - 2$ and this was proved simultaneously and independently by GRÜNBAUM [3], HEPPES [4] and STRASZEWICZ [5] (all using similar methods). I am going to prove

$$(2) \quad c_1 \cdot n^{4/3} < G_3(n) < c_2 \cdot n^{5/3}.$$

Perhaps $G_k(n) < n^{4/3+\epsilon}$ holds for all $n > n(\epsilon)$.

One could have expected that $G_k(n) = o(n^2)$ and $g_k(n) < c_k \cdot n$ for every k . In 1955 LENZ showed that this is not so. In fact LENZ showed that (LENZ's result is unpublished)

$$(3) \quad g^*(n) \cong \frac{n^2}{4}.$$

The proof of LENZ is very **simple**. Put $s = \left\lfloor \frac{n}{2} \right\rfloor$ and consider the following n points in four-dimensional space:

$$(x_i, y_i, 0, 0), \quad 1 \leq i \leq s, \quad (0, 0, x_j, y_j), \quad s+1 \leq j \leq n$$

where $0 < x_i, x_j, y_i, y_j < \frac{1}{\sqrt{2}}$, $x_i^2 + y_i^2 = \frac{1}{2}$, $x_j^2 + y_j^2 = \frac{1}{2}$. Clearly all the

is $(n - s) = \left\lfloor \frac{n}{2} \right\rfloor$ distances between the points $(x_i, y_i, 0, 0)$ and $(0, 0, x_j, y_j)$ is 1 (and 1 is the diameter of the set $(x_i, y_i, 0, 0); (0, 0, x_j, y_j)$).

By a slight modification of this method LENZ in fact proved that $g_k(n) \geq \frac{n^2}{4} + c_3 n$ for a certain $c_3 > 0$. LENZ then asked: what is the limit of $g_k(n)/n^2$ as $n \rightarrow \infty$? In this note I am going to prove the following

Theorem. For every $k \geq 4$

$$\lim_{n \rightarrow \infty} g_k(n)/n^2 = \lim_{n \rightarrow \infty} G_k(n)/n^2 = \frac{1}{2} - \frac{1}{2 \left\lfloor \frac{k}{2} \right\rfloor}$$

Clearly $g(n) \leq G(n)$ and $g_k(n) \leq g_{k+1}(n)$, $G(n) \leq G_{k+1}(n)$. Thus to prove our Theorem it will suffice to show that for every $l \geq 2$

$$(4) \quad \lim_{n \rightarrow \infty} g_{2l}(n)/n^2 \geq \frac{1}{2} - \frac{1}{2l}$$

and

$$(5) \quad \lim_{n \rightarrow \infty} G_{2l+1}(n)/n^2 \leq \frac{1}{2} - \frac{1}{2l}.$$

The proof of (4) is trivial generalization of the proof of LENZ. For each $l \leq t \leq l$ denote by I_t the group of $\left\lfloor \frac{n}{l} \right\rfloor$ points whose first $2t - 2$ coordinates are 0 the $2t - 1$ -th and $2t$ -th coordinates are x_i, y_i , $1 \leq i \leq \left\lfloor \frac{n}{l} \right\rfloor$, $x_i, y_i > 0, x_i^2 + y_i^2 = \frac{1}{2}$ and the remaining $2l - 2t$ coordinates are 0. Clearly for any $t_1 \neq t_2$ the distance between any two points of I_{t_1} and I_{t_2} is 1 and the set $\bigcup_{1 \leq t \leq l} I_t$ has diameter 1. Thus

$$g_{2l}(n) \geq \binom{l}{2} \left[\frac{n}{l} \right]^2 = \frac{n^2}{2} \left(1 - \frac{1}{l} \right) + O(n)$$

which clearly implies (4).

Next we prove (5). If (5) is not true then there exists an $\varepsilon > 0$ so that for a certain $l \geq 2$ and infinitely many n_s

$$G_{2l+1}(n_s) > \left(\frac{1}{2} - \frac{1}{2l} + \varepsilon \right) n_s^2 = A(n_s)$$

In other words there exists a set $P_{n_s}^{(2l+1)}$ in $2l + 1$ dimensional space and a distance r which occurs among at least $A(n_s)$ pairs of points of $P_{n_s}^{(2l+1)}$. Connect any two points of $P_{n_s}^{(2l+1)}$ whose distance is r . Thus we obtain a graph

of n_s vertices and $A(n_s)$ edges. By a theorem of A. H. STONE and myself [6] this graph contains for sufficiently large $n_s = n_s(\varepsilon)$ a subgraph of $3(l+1)$ vertices $x_i^{(l)}$ $1 \leq i \leq 3, 1 \leq l \leq l+1$ so that any two vertices $x_{i_1}^{(l)}$ and $x_{i_2}^{(l)}$ are connected by an edge if $t_{i_1} \neq t_{i_2}$ (in other words the distance between $x_{i_1}^{(l)}$ and $x_{i_2}^{(l)}$ is r if $t_{i_1} \neq t_{i_2}$). But, then a simple geometrical argument shows that the $l+1$ planes $(x_1^{(l)}, x_2^{(l)}, x_3^{(l)})$ $1 \leq l \leq l+1$ must be mutually perpendicular, which implies that the dimension of the space spanned by the $x_i^{(l)}$ is at least $2l+2$. This contradiction proves (5) and thus the proof of our Theorem is complete.

By a sharpening which I recently obtained of the result of STONE and myself I can prove

$$(6) \quad G_k(n) < \left(\frac{1}{2} - \frac{1}{2 \lfloor \frac{k}{2} \rfloor} \right) n^2 + O(n^{2-\varepsilon_k})$$

where $\varepsilon_k \rightarrow 0$ as $k \rightarrow \infty$. I do not know how close (6) is to the true order of magnitude of $G_k(n)$. Perhaps the result of LENZ

$$(7) \quad G_k(n) > \left(\frac{1}{2} - \frac{1}{2 \lfloor \frac{l}{k} \rfloor} \right) n^2 + c_k n$$

gives the right order of magnitude.

Now we are going to prove (2). First we prove the upper estimate. Let x_1, x_2, \dots, x_n be n points in three dimensional space, assume that there are α_i points at distance r from x_i . Clearly to any three points $x_{j_1}, x_{j_2}, x_{j_3}$ there can be at most two points x_i at distance r . Thus since the total number of triplets $(x_{j_1}, x_{j_2}, x_{j_3})$ is $\binom{n}{3}$ a simple argument gives

$$\sum_{i=1}^n \binom{\alpha_i}{3} \leq 2 \binom{n}{3}$$

or

$$(8) \quad \sum_{i=1}^n \alpha_i^3 < c_4 n^3.$$

If $\sum_{i=1}^n \alpha_i^3$ is given $\sum_{i=1}^n \alpha_i$ is maximal if all the α_i are equal. Thus (8) implies

$$\sum_{i=1}^n \alpha_i < c_2 n^{3/2}$$

which proves the upper bound in (2).

¹ The theorem in question states as follows: To every $\varepsilon, r \geq 2$ and l there exists an $n_0(\varepsilon, r, l)$ so that if $n > n_0(\varepsilon, r, l)$ and G_n is a graph of n vertices and more than $n^2 \left(\frac{1}{2} - \frac{1}{2(r-1)} + \varepsilon \right)$ edges then G_n contains rl vertices $x_i^{(l)}$ $1 \leq i \leq l, 1 \leq i_1 \leq r$ so that for every $i_1 \neq i_2, x_{i_1}^{(l)}$ and $x_{i_2}^{(l)}$ are connected by an edge for every $1 \leq i_1, i_2 \leq l$.

To prove the lower bound in (2) consider the points (x_i, y_i, z_i) of integer coordinates $0 \leq x_i, y_i, z_i \leq [n^{1/3}]$. Clearly the number of these points is less than n but is greater than $n(1 - \varepsilon)$. The square of the distance between two of these points is of the form

$$(9) \quad u^2 + v^2 + w^2, \quad 0 \leq u, v, w \leq n^{1/3}$$

The numbers (9) are all less than or equal to $3n^{2/3}$ and since there are more than $\binom{n(1 - \varepsilon)}{2}$ such distances, clearly for some n the same distance must occur at least $1/7n^{1/3}$ times, which completes the proof of (2). From deep number theoretic results it follows that for suitable r the same distance occurs more than $c_3 n^{1/3} \log \log n$ times and this is the **best** lower bound I can get for $G_3(n)$ at the present time.

(Received December 18, 1959.)

REFERENCES

- [1] Aufgabe 167. *Jahresbericht der Deutschen Math. Vereinigung* **43** (1934) 114.
 [2] ERDŐS, P.: „On sets of distances of n points.” *Amer. Math. Monthly* **53** (1946) 248, 250.
 [3] GRÜNBAUM, B.: “A proof of Vázsonyi’s conjecture,” Bull. *Research Council of Israel* **6A** (1956) 77-78.
 [4] HEPPES, A.: „Beweis einer Vermutung von A. Vázsonyi.” *Acta Math Acad. Sci. Hung.* **7** (1957) 463-466.
 [5] STRASZEWICZ, S.: “Sur un problème géométrique de P. Erdős.” *Bull. Acad. Pol. Sci. Cl. III* **5** (1957) 39-40.
 [6] ERDŐS, P. and STONE, A. H.: “On the structure of linear graphs.” *Bull. Amer. Math. Soc.* **52** (1946) 1087-1091.

О РАССТОЯНИЯХ МЕЖДУ n ТОЧКАМИ ЭВКЛИДОВА ПРОСТРАНСТВА

P. ERDŐS

Резюме

Пусть $P_n^{(k)}$ есть множество, состоящее из n точек k -мерного пространства диаметра которого равен 1. Обозначим через $g_k(n, r)$ максимальное число пар точек (x_i, x_j) для которых расстояния x_i и x_j равно r .

$$G_k(n) = \max_{(r)} g_k(n, r); \quad g_k(n) = g_k(n, 1).$$

Раньше автор доказал, что

$$n^{1-c_1/\log \log n} < G_2(n) < n^{3/2}.$$

Было известно, что $g_2(n) = n$. GRÜNBAUM, HEPPES и STRASZEWICZ доказали гипотезу VAZSONYI, согласно которой $g_3(n) = 2n - 2$. LENZ доказал, что

$$g_3(n) > \frac{n^2}{4} + c_2 n.$$

В настоящей статье автор доказывает, что

$$c_3 n^{4/3} < G_3(n) < c_4 n^{5/3}$$

и если $k \geq 4$, то

$$\lim_{n \rightarrow \infty} g_k(n)/n^2 = \lim_{n \rightarrow \infty} G_k(n)/n^2 = \frac{1}{2} - \frac{1}{2 \left\lfloor \frac{k}{2} \right\rfloor}.$$